

Formalization of Ownership Transfer in Universe Types

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Abstract

Ownership is a powerful concept to structure the object store and to control aliasing and modifications of objects. This paper presents an ownership type system for a Java-like programming language. Like our earlier Universe type system, Universe Types with Transfer (UTT) enforce the owner-as-modifier discipline. This discipline does not restrict aliasing, but requires modifications of an object to be initiated by its owner. This allows owner objects to control state changes of owned objects, for instance, to maintain invariants. UTT combines ownership type checking with a modular static analysis to control references to transferable objects. UTT is very flexible because it permits temporary aliases, even across certain method calls. Nevertheless, it guarantees statically that a cluster of objects is externally-unique when it is transferred and, thus, that ownership transfer is type safe. UTT provides the same encapsulation as Universe Types and requires only negligible annotation overhead.

Contents

1 Syntax	3
2 Auxiliary functions and definitions	4
2.1 General constructions	4
2.2 Clusters	5
2.3 Ownership type modifiers	5
2.4 Type combinators	5
2.5 Subtype relation	6

2.6	Universe types's function	6
2.7	Lookup functions	6
2.8	Unusable variables	8
3	Unusable set's generation rules	9
4	Type rules	10
5	Runtime Model	12
5.1	Heap Model	12
5.2	Operations on Heap and Objects	13
5.3	Operational semantic	14
6	Type safety	17
7	Type safety proof	19
7.1	The main theorem	19
7.2	Auxiliary lemmas	26
7.3	heap's class well-formedness lemmas	32
7.4	heap's well-formedness lemmas	33
7.5	Frames stack's well-formedness lemmas	35
7.6	Uniqueness's lemmas	40
7.7	Global invariant's lemmas	43
7.8	Ownership tree's lemmas	47
7.9	This owners's lemmas	48
7.10	Owner as modifier's lemmas	51

1 Syntax

$CL ::= \text{class } C \text{ extends } D \{ \overline{Cl}; \overline{T} \overline{f}; \overline{M} \}$

$M ::= T \text{ } mt(T \text{ } x) \text{ } locVar(\overline{T} \overline{y}) \{ e; \}$

$e ::=$

- $x = y$
- $| x = y.f$
- $| x.f := y$
- $| x = y.mt(z)$
- $| x = new \text{ } T()$
- $| x = (T)y$
- $| e_1; e_2$
- $| x = release(y)$
- $| x = capture \langle m \rangle (y)$

$T ::= m \text{ } C$

$m ::= \text{any} \mid \text{peer} \mid \text{rep} \langle Cl \rangle \mid \text{uniq}$

Notation:

- $x, y, z \in \text{VarID}$ ranges over local variables and formal parameters, including `this`
- $C \in \text{ClassID}$ ranges over classes' names
- $mt \in \text{MethodID}$ ranges over methods' names
- $f \in \text{FieldID}$ ranges over fields' names
- $T \in \text{Type}$ ranges over Universe types
- $m \in \text{MOD}$ ranges over Universe types modifiers
- $Cl \in \text{Clt}$ ranges over clusters

Note: In purpose of simplification we:

- forbidden nested expression. Any program can be transformed to this form via introducing additional local variables which represent temporal values.

- suppose that any method has only one parameter. It is not hard to generalize in case of arbitrary parameters number.
- suppose that all local variables are declared at the begin of a method. It is not hard transform a program to this form via collection variables declaration in the method body.
- to return result of a method invocation we use special local variable `res`.
- All classes' fields and local variable have unique names.
- All input parameters have fixed name p .
- If methods have equal names then they have equal signatures.

2 Auxiliary functions and definitions

2.1 General constructions

In purpose of abbreviation we use next constructions:

$$\begin{aligned}
 x = & \quad \text{if}(b_1) \quad \text{then } t_1 \\
 & \quad \text{else} \quad \text{if}(b_2) \quad \text{then } t_2 \\
 & \quad \quad \ldots \\
 & \quad \quad \quad \text{else} \quad \text{if}(b_n) \quad \text{then } t_n \\
 & \quad \quad \quad \text{else} \quad t_{n+1}
 \end{aligned}
 \Leftrightarrow \bigwedge_{i=1}^{n+1} \left(\bigwedge_{j=1}^{i-1} \neg b_j \wedge b_i \Rightarrow x = t_i \right)$$

where $b_{n+1} = \top$

We use next operations on sequences, and mappings:

- $\epsilon = \langle \rangle$
- $\langle a_1, \dots, a_n \rangle \circ \langle b_1, \dots, b_m \rangle = \langle a_1, \dots, a_n, b_1, \dots, b_m \rangle$
- $\langle a_1, \dots, a_n \rangle \downarrow_i = \text{if}(i \in [1..n]) \text{ then } a_i \text{ else undefined}$
- $M[a \mapsto b] = (M \setminus \{\langle a, _ \rangle\}) \cup \{\langle a, b \rangle\}$
- $\text{dom}(\{\langle a_1, b_1 \rangle, \dots, \langle a_n, b_n \rangle\}) = \{a_1, \dots, a_n\}$
- $\text{rng}(\{\langle a_1, b_1 \rangle, \dots, \langle a_n, b_n \rangle\}) = \{b_1, \dots, b_n\}$

2.2 Clusters

Clt is an unrestricted set which contains all clusters.

$$\text{Clt} = \text{CltRep} \cup \text{CltFree} \quad \text{CltRep} \cap \text{CltFree} = \emptyset$$

CltRep is an unrestricted set which contains all clusters owned by some object.

CltFree is an unrestricted set which contains all unowned clusters.

We have a special cluster Cl_{this} which contains object's representation.

$$Cl_{\text{this}} \in \text{CltRep}$$

To be able generate new unique cluster we assume that CltFree has next structure: $\text{CltFree} = \{\text{CF}_i | i \in \mathbb{N}\}$

CF_i satisfies the next axiom: $\forall i, j \ i \neq j \Rightarrow \text{CF}_i \neq \text{CF}_j$

2.3 Ownership type modifiers

A user can use next type modediers: `any`, `peer`, `rep` $\langle Cl \rangle$, `uniq`. Also we use type modifier `this`. A user can't use it directly but it can be deduced during the type checking procedure invocation. `this` is type modifier which corresponds to the `this` reference.

We denote meta-type modifiers like

$$\text{MOD} = \{\text{any}, \text{peer}, \text{rep } \langle Cl \rangle, \text{uniq}, \text{this} \mid Cl \in \text{Clt}\}$$

2.4 Type combinators

We define a type combinator \triangleright_U . We use it for expressing field access and methods invocation.

$$\begin{array}{llll} \text{peer} & \triangleright_U & \text{peer} & \rightarrow \text{peer} \\ \text{rep } \langle Cl \rangle & \triangleright_U & \text{peer} & \rightarrow \text{rep } \langle Cl \rangle \\ \text{this} & \triangleright_U & x & \rightarrow x \\ x & \triangleright_U & \text{uniq} & \rightarrow \text{uniq} \\ x & \triangleright_U & y & \rightarrow \text{any} \end{array}$$

We define a type combinator \triangleright_U in the next way:

$$(m_1 \ C_1) \triangleright_U (m_2 \ C_2) = (m_1 \triangleright_U m_2) \ C_2$$

2.5 Subtype relation

We define subtype relation on the Java classes in the next way:

$$C \sqsubseteq C \quad \frac{C \sqsubseteq D \quad D \sqsubseteq E}{C \sqsubseteq E} \quad \frac{\text{class } C \text{ extends } D \{ \dots \}}{C \sqsubseteq D}$$

We define subtype relation on the universe type's modifiers in the next way:

$$\frac{m = m' \vee m' = \text{any}}{m \leq m'}$$

We define subtype relation on the universe types in the next way:

$$\frac{m \leq m' \quad C \sqsubseteq C'}{m \ C \leq m' \ C'}$$

2.6 Universe types's function

To access universe types' components we use next functions:

$$\frac{T = m \ C}{\text{class}(T) = C} \quad \frac{T = m \ C}{\text{mod}(T) = m}$$

In purpose of brevity we allow use universe type in any expression in which we can use universe type modifiers. For example if we expect that expression $F[_]$ get as parameter type modifier then $F[T] \stackrel{\text{def}}{=} F[\text{mod}(T)]$

2.7 Lookup functions

- The function $\text{fields}(C)$ yields the identifiers of all fields that are declared in or inherited by class C .

$$\frac{}{\text{fields}(Object) = \epsilon} \quad \frac{\text{class } C \text{ extends } D \{ _ ; \overline{T}f; _ \}}{\text{fields}(C) = \overline{f} \circ \text{fields}(D)}$$

- The function $\text{fType}(C, f)$ yields the type of field f as declared in class C . The result is `undefined` if f is not declared in C . Since identifiers

are assumed to be globally unique, there is only one declaration for each field identifier.

$$\frac{\text{class } C \text{ extends } \dots \{ \dots ; \dots T f \dots ; \dots \}}{\mathbf{fType}(C, f) = T}$$

- The function $\mathbf{mType}(C, mt)$ yields the signature of method mt as declared in class C . The result is `undefined` if mt is not declared in C . We do not allow overloading of methods; therefore, the method identifier is sufficient to uniquely identify a method.

$$\frac{\text{class } C \text{ extends } \dots \{ \dots ; \dots ; \dots T_r \text{ } mt(T_p p) \dots \}}{\mathbf{mType}(C, mt) = T_p \rightarrow T_r}$$

- The function $\mathbf{mBody}(C, mt)$ yields a mt 's body expression.

$$\frac{\text{class } C \text{ extends } \dots \{ \dots ; \dots ; \dots \text{ } mt(\dots) \dots \{ e; \} \dots \}}{\mathbf{mBody}(C, mt) = e}$$

$$\frac{\text{class } C \text{ extends } D \{ \text{no method } mt \}}{\mathbf{mBody}(C, mt) = \mathbf{mBody}(D, mt)}$$

- Method mt respects the rules for overriding if it does not override a method or if all overridden methods have the identical signatures.

$$\frac{\forall C' : C' \leq C \Rightarrow \\ \mathbf{mType}(C', mt) \text{ undefined} \vee \mathbf{mType}(C) = \mathbf{mType}(C')}{\mathbf{override}(C, mt)}$$

- The function \mathbf{mLoc} yields the names of local variables as declared in class C :

$$\frac{\text{class } C \text{ extends } \dots \{ \dots ; \dots ; \dots \text{ } mt(\dots) locVar(\overline{T} \bar{y}) \{ \dots \} \dots \}}{\mathbf{mLoc}(C, mt) = \bar{y}}$$

- The function \mathbf{Clt}_C yields the clusters of class C :

$$\frac{\text{class } C \text{ extends } \dots \{ \overline{Cl}; \dots ; \dots \}}{\mathbf{Clt}_C = \overline{Cl}}$$

- The function \mathbf{CltAll}_C yields the clusters of class C and its superclasses:

$$\frac{}{\mathbf{CltAll}_{Object} = \emptyset} \quad \frac{\text{class } C \text{ extends } D \dots}{\mathbf{CltAll}_C = \mathbf{Clt}_C \cup \mathbf{CltAll}_D}$$

2.8 Unusable variables

Note: we consider unusable variables set for fixed:

- Class C
- Method mt
- Statical environment Γ

To represent data flow analyze information we use special set which contains all inaccessible (unusable) variables. It is subset of $\mathbf{Var} = \mathbf{Loc} \cup \mathbf{Field}$ where:

- $\mathbf{Loc} = \{v \in \mathbf{dom}(\Gamma) \mid \Gamma(v) \in \{\mathbf{rep}\langle Cl \rangle, \mathbf{uniq}\}\}$
- $\mathbf{Field} = \{f \in \mathbf{fields}(C, mt) \mid \mathbf{fType}(C, f) = \mathbf{rep}\langle Cl \rangle\}$

To change am unusable set we use next operations:

- $\mathbf{consumeAliases}(U, v) = U \cup \{x \in \mathbf{Var} \mid \Gamma(v) = \Gamma(x)\} \cup \{f \in \mathbf{fields}(C, mt) \mid \Gamma(v) = \mathbf{fType}(C, f)\}$
- $\mathbf{consumeUniq}(U, T_{res}, v) = \begin{cases} \text{if } (T_{res} \neq \text{any} \wedge \Gamma(v) = \mathbf{uniq}) \text{ then } U \cup \{v\} \text{ else } U \end{cases}$
- $\mathbf{consumeLocals}(U) = U \cup \{v \in \mathbf{Loc} \mid \Gamma(v) = \mathbf{rep}\langle Cl_f \rangle \wedge f \in \mathbf{Field}\}$

3 Unusable set's generation rules

In this section we express change of unusable set during expression type checking or evaluation. We use $\Gamma \in \text{Env} = \overline{\text{VarID Type}}$ for the declaration environment, which maps formal parameters to their types. A unusable set's generation rules has the form $\Gamma; U \vdash_{DE} e : U'$ and expresses that after expression e type checking or evaluation the unusable set change from U to U' .

In some situations we need to know an unusable set before method invocation. **UU-PRE-INVK** rule generate such unusable set.

$$\begin{array}{c}
\frac{U' = \text{consumeUniq}(U, \Gamma(x), y) \setminus \{x\}}{\Gamma; U \vdash_{DE} x = y : U'} \quad [\text{UU-ASSIGN}] \\
\\
\frac{}{\Gamma; U \vdash_{DE} x = y.f : U \setminus \{x\}} \quad [\text{UU-FIELD-READ}] \\
\\
\frac{U' = \text{if}(y = \text{this}) \\ \text{then } U \setminus \{f\} \text{ else } U}{\Gamma; U \vdash_{DE} y.f = x : U'} \quad [\text{UU-FIELD-WRITE}] \\
\\
\frac{\text{mType}(\Gamma(\text{class}(y)), mt) = T_p \rightarrow T_r \\ U_1 = \text{consumeUniq}(U, T_p, z) \\ U_2 = \text{if}(\Gamma(y) \notin \{\text{peer, this}\}) \text{ then } U_1 \\ \text{else consumeLocals}(\text{consumeUniq}(U_1, T_p, z))}{\Gamma; U \vdash_{DE} y.mt(z) : U'} \quad [\text{UU-PRE-INVK}] \\
\\
\frac{\Gamma; U \vdash DEy.mt(z) : U'}{\Gamma; U \vdash_{DE} x = y.mt(z) : U' \setminus \{x\}} \quad [\text{UU-INVK}] \\
\\
\frac{}{\Gamma; U \vdash_{DE} x = \text{new } T() : U \setminus \{x\}} \quad [\text{UU-NEW}] \\
\\
\frac{U' = \text{consumeUniq}(U, \Gamma(x), y) \setminus \{x\}}{\Gamma; U \vdash_{DE} x = (T)y : U'} \quad [\text{UU-CAST}]
\end{array}$$

$$\begin{array}{c}
\frac{\Gamma; U \vdash_{DE} e_1 : U_1 \quad \Gamma; U_1 \vdash_{DE} e_2 : U_2}{\Gamma; U \vdash_{DE} e_1; e_2 : U_2} \quad [\text{UU-COMP}] \\
\frac{U' = \text{consumeAliases}(U, y) \setminus \{x\}}{\Gamma; U \vdash_{DE} x = \text{release}(y) : U'} \quad [\text{UU-RELEASE}] \\
\frac{U' = U \cup \{y\} \setminus \{x\}}{\Gamma; U \vdash_{DE} x = \text{capture} \langle m \rangle(y) : U'} \quad [\text{UU-CAPTURE}]
\end{array}$$

4 Type rules

A judgment has the form $\Gamma; U \vdash e$ and expresses that expression e is well-typed in a declaration environment Γ . U is an unusable set before the expression evaluation.

We use an auxiliary predicate $\text{RepInDom} : \text{MOD} \times \mathcal{P}(\text{Clt})$ to check if the type's domain belongs to the proper domains set or not.

$$\text{RepInDom}(m, CltSet) \stackrel{\text{def}}{=} m = \text{rep} \langle Cl \rangle \Rightarrow Cl \in CltSet$$

$$\begin{array}{l}
\text{IsWritable} : \text{MOD} \rightarrow \text{Bool} \\
\text{IsWritable}(m) = m \in \{\text{peer}, \text{this}, \text{rep} \langle Cl \rangle \mid Cl \in \text{Clt}\}
\end{array}$$

The function $\Gamma_{C, mt}$ yields the static environment of method mt as declared in class C :

$$\frac{\text{class } C \text{ extends } _ \{ _ ; _ ; \dots T_r \ mt(T_p \ x) \ locVar(\overline{T \ y}) \ \{ _ \} \dots \}}{\Gamma_{C, mt} = \{\langle p, T_p \rangle, \ \overline{\langle y, T \rangle}, \ \langle \text{this}, \text{this } C \rangle, \ \langle \text{res}, T_r \rangle\}}$$

$$\frac{y \notin U}{\Gamma(y) \leq \Gamma(x)} \quad \text{[T-ASSIGN]}$$

$$\frac{\begin{array}{c} T_f = \mathbf{fType}(\mathbf{class}(\Gamma(y)), f) \\ \Gamma(y) \triangleright_U T_f \leq \Gamma(x) \\ y \notin U \\ y = \mathbf{this} \Rightarrow f \notin U \end{array}}{\Gamma; U \vdash x = y.f} \quad \text{[T-FIELD-READ]}$$

$$\frac{\begin{array}{c} T_f = \mathbf{fType}(\mathbf{class}(\Gamma(y)), f) \\ \Gamma(x) \leq \Gamma(y) \triangleright_U T_f \\ \mathbf{IsWritable}(\Gamma(y)) \\ T_f = \mathbf{rep}\langle Cl \rangle \Rightarrow \Gamma(y) = \mathbf{this} \\ x, y \notin U \end{array}}{\Gamma; U \vdash y.f = x} \quad \text{[T-FIELD-WRITE]}$$

$$\frac{\begin{array}{c} \mathbf{mType}(C, mt) = T_p \rightarrow T_r \\ \Gamma(z) \leq \Gamma(y) \triangleright_U T_p \\ \Gamma(y) \triangleright_U T_r \leq \Gamma(x) \\ T_p = \mathbf{rep}\langle Cl \rangle \Rightarrow \Gamma(y) = \mathbf{this} \\ y, z \notin U \\ \Gamma; U \vdash_{DE} y.mt(z) : U' \\ \mathbf{Field} \cap U' = \emptyset \end{array}}{\Gamma; U \vdash x = y.mt(z)} \quad \text{[T-INVK]}$$

$$\frac{\begin{array}{c} \mathbf{IsWritable}(T) \\ C = \mathbf{class}(\Gamma(\mathbf{this})) \\ \mathbf{RepInDom}(T, \mathbf{CltAll}_C) \\ T \leq \Gamma(x) \end{array}}{\Gamma; U \vdash x = \mathbf{new}\ T()} \quad \text{[T-NEW]}$$

$$\frac{\begin{array}{c} T \leq \Gamma(x) \\ C = \mathbf{class}(\Gamma(\mathbf{this})) \\ \mathbf{RepInDom}(T, \mathbf{CltAll}_C) \\ m_y = \mathbf{mod}(\Gamma(y)) \\ m = \mathbf{mod}(T) \\ m_y \leq m \vee (m_y = \mathbf{any} \wedge \mathbf{IsWritable}(\Gamma(m))) \\ y \notin U \end{array}}{\Gamma; U \vdash x = (T)y} \quad \text{[T-CAST]}$$

$$\begin{array}{c}
\frac{\Gamma; U \vdash e_1 \quad \Gamma; U \vdash_{DE} e_1 : U' \quad \Gamma; U' \vdash e_2}{\Gamma; U \vdash e_1; e_2} \quad [\text{T-COMP}] \\[10pt]
\frac{\Gamma(x) = \text{uniq} \quad \Gamma(y) = \text{rep} \langle Cl \rangle \quad Cl \neq Cl_{\text{this}} \quad \text{class}(\Gamma(y)) \sqsubseteq \text{class}(\Gamma(x)) \quad y \notin U}{\Gamma; U \vdash x = \text{release}(y)} \quad [\text{T-RELEASE}] \\[10pt]
\frac{\text{IsWritable}(m) \quad C = \text{class}(\Gamma(\text{this})) \quad \text{RepInDom}(m, \text{CltAll}_C) \quad \Gamma(x) = m \quad \Gamma(y) = \text{uniq} \quad \text{class}(\Gamma(y)) \sqsubseteq \text{class}(\Gamma(x)) \quad y \notin U}{\Gamma; U \vdash x = \text{capture} \langle m \rangle (y)} \quad [\text{T-CAPTURE}] \\[10pt]
\frac{\text{RepInDom}(\overline{T} \circ T_r \circ T_p, \text{CltAll}_C) \quad \Gamma_{C, mt}; \emptyset \vdash e \quad \Gamma_{C, mt}; \emptyset \vdash_{DE} e : U \quad (\text{Field} \cup \{\text{res}\}) \cap U = \emptyset}{T_r \ mt(T_p \ x) \ locVar(\overline{T} \ y) \ \{e; \} \in C} \quad [\text{WF-METHOD}] \\[10pt]
\frac{\overline{M} \in C \quad \text{RepInDom}(\overline{f}, \text{Clt}_C) \quad \overline{f} \neq \text{uniq}}{\text{class } C\{\overline{Cl}; \ \overline{T} \ \overline{f}; \ \overline{M}\}} \quad [\text{WF-CLASS}]
\end{array}$$

5 Runtime Model

5.1 Heap Model

Fig. 1 summarizes our heap model. To distinguish sorts of the runtime model from their static counterparts, we use the prefix r .

$h \in \text{Heap}$	$= \text{Addr!} \rightarrow \text{Obj}$
$\iota! \in \text{Addr!}$	$= \text{Set of Addresses}$
$\iota \in \text{Addr}$	$= \text{Addr!} \cup \{\text{null}^r\}$
$o \in \text{Obj}$	$= \text{ClassID} \times \text{Owners} \times \text{Fields}$
$\text{Fds} \in \text{Fields}$	$= \text{FieldID} \rightarrow \text{Addr}$
$\text{ow} \in \text{Owners}$	$= \text{Clt} \times \text{Addr}$
$\Gamma^r \in \text{Env}^r$	$= \overline{\text{VarID} \text{ Addr}}$
$\text{Fr} \in \text{Frames}$	$= \text{Env}^r \times \text{Env} \times \mathcal{P}(\text{VarID})$
$\text{FrSt} \in \text{FramesStack}$	$= \overline{\text{Frames}}$
$\sigma \in \text{States}$	$= \text{Heap} \times \text{FramesStack}$
$\sigma_{pr} \in \text{States}_{pr}$	$= \text{Heap} \times \text{Env}^r$

Figure 1: Definitions for the heap model.

A heap (sort Heap) maps addresses to objects. The set of addresses (Addr) contains the special null-reference null^r . An object (Obj) consist of its runtime type and a mapping from field identifiers to the addresses stored in the fields.

5.2 Operations on Heap and Objects

$$\begin{aligned}
 \cdot[\dots := \cdot] &:: \text{Heap} \times \text{Addr!} \times \text{FieldID} \times \text{Addr} \rightarrow \text{Heap} \\
 h[\iota.f := \iota'] &= h[\iota \mapsto \langle h(\iota) \downarrow_1, h(\iota) \downarrow_2, h(\iota) \downarrow_3 [f \mapsto \iota'] \rangle] \\
 \cdot(\dots) &:: \text{Heap} \times \text{Addr!} \times \text{FieldID} \rightarrow \text{Addr} \\
 h(\iota.f) &= h(\iota) \downarrow_3 (f) \\
 \text{class}^r &:: \text{Heap} \times \text{Addr!} \rightarrow \text{ClassID} \\
 \text{class}^r(h, \iota) &= h(\iota) \downarrow_1 \\
 \text{owner} &:: \text{Heap} \times \text{Addr!} \rightarrow \text{Owners} \\
 \text{owner}(h, \iota) &= h(\iota) \downarrow_2 \\
 \text{owners} &:: \text{Heap} \times \text{Addr!} \rightarrow \mathcal{P}(\text{Owners}) \\
 \text{owners}(h, \iota) &\in \text{owners}(h, \iota) \\
 \iota' \in \text{owners}(h, \iota) \wedge \iota' \downarrow_2 \neq \text{null}^r &\Rightarrow \text{owner}(h, \iota' \downarrow_2) \in \text{owners}(h, \iota)
 \end{aligned}$$

$$\begin{array}{lcl} \text{transfer} :: \text{Heap} \times \text{Addr!} \times \text{Owners} \rightarrow \text{Heap} \\ \text{transfer}(\text{h}, \iota, \text{ow}) & = & \text{h}[\iota \mapsto \langle \text{ow}, \text{h}(\iota) \downarrow_2, \text{h}(\iota) \downarrow_3 \rangle] \end{array}$$

$$\begin{array}{lcl} \text{transferCl} :: \text{Heap} \times \text{Addr} \times \text{Owners} \rightarrow \text{Heap} \\ \text{transferCl}(\text{h}, \iota, \text{ow}) & = & \text{if } (\iota = \text{null}^r) \text{ then h} \\ & & \text{else transfer}(\dots \text{transfer}(\text{transfer}(\text{h}, \iota_1, \text{ow}), \iota_2, \text{ow}) \dots, \iota_n, \text{ow}) \\ & & \text{where } \{\iota_1, \iota_2, \dots, \iota_n\} = \{\iota' \mid \text{owner}(\text{h}, \iota') = \text{owner}(\text{h}, \iota)\} \end{array}$$

$$\begin{array}{lcl} \text{newClt} :: \text{Heap} \rightarrow \text{CltFree} \\ Cl = \text{newClt}(\text{h}) & \Rightarrow & Cl \notin \text{owner}(\text{h}, \text{dom}(\text{h})) \downarrow_1 \end{array}$$

$$\begin{array}{lcl} \text{mod20w} :: \text{Heap} \times \text{MOD} \times \text{Addr!} \rightarrow \text{Owners} \\ \text{mod20w}(\text{h}, m, \iota) = \text{if } (m = \text{rep} \langle Cl \rangle) \text{ then } \langle Cl, \iota \rangle \\ \text{else if } (m = \text{peer}) \text{ then } \text{owner}(\text{h}, \iota) \\ \text{else if } (m = \text{uniq}) \text{ then } \langle \text{newClt}(\text{h}), \text{null}^r \rangle \\ \text{else } \text{undefined} \end{array}$$

$$\begin{array}{lcl} \text{new} :: \text{Heap} \times \text{ClassID} \times \text{Owners} \rightarrow \text{Addr!} \times \text{Heap} \\ \text{new}(\text{h}, C, \text{ow}) & = & \langle \iota, \text{h}[\iota \mapsto \langle C, \text{ow}, \text{FdS} \rangle] \rangle \\ \text{where } \iota \notin \text{dom}(\text{h}) \wedge \text{FdS}(\text{fields}(C)) = \text{null}^r \end{array}$$

$$\begin{array}{lcl} \text{prj} :: \text{States} \rightarrow \text{States}_{pr} \\ \text{prj}(\langle \text{h}, \langle \Gamma^r, \Gamma, U \rangle \circ \text{FrSt} \rangle) & = & \langle \text{h}, \Gamma^r \rangle \end{array}$$

$$\begin{array}{l} \text{RefNum} :: \text{Heap} \times \text{Frames} \times \text{Owners} \times \text{MOD} \rightarrow \mathbb{N} \\ \text{RefNum}(\text{h}, \langle \Gamma^r, \Gamma, U \rangle, \text{ow}, m) = \\ | \{x \mid \text{mod}(\Gamma(x)) = m \wedge \text{owner}(\text{h}, \Gamma^r(x)) = \text{ow} \wedge x \notin U\} | \end{array}$$

$$\begin{array}{l} \text{RefNum} :: \text{Heap} \times \text{FramesStack} \times \text{Owners} \times \text{MOD} \rightarrow \mathbb{N} \\ \text{RefNum}(\text{h}, \overline{\text{Fr}}, \text{ow}, m) = \sum_{\text{Fr} \in \overline{\text{Fr}}} \text{RefNum}(\text{h}, \text{Fr}, \text{ow}, m) \end{array}$$

5.3 Operational semantic

A judgment has the form $\text{h}; \text{FrSt}; e \rightsquigarrow \text{h}'; \Gamma^r$ and expresses evaluation of an expression e .

$\frac{\Gamma^r' = \Gamma^r[x \mapsto \Gamma^r(y)]}{\mathbf{h}; \langle \Gamma^r, \Gamma, U \rangle \circ \text{FrSt}; x = y \rightsquigarrow \mathbf{h}; \Gamma^r}$	[OS-ASSIGN]
$\frac{\Gamma^r(y) \neq \text{null}^r \quad \Gamma^r' = \Gamma^r[x \mapsto \mathbf{h}(\Gamma^r(y).f)]}{\mathbf{h}; \langle \Gamma^r, \Gamma, U \rangle \circ \text{FrSt}; x = y.f \rightsquigarrow \mathbf{h}; \Gamma^r}$	[OS-FIELD-READ]
$\frac{\Gamma^r(x) \neq \text{null}^r \quad \mathbf{h}' = \mathbf{h}[\Gamma^r(x).f := \Gamma^r(y)]}{\mathbf{h}; \langle \Gamma^r, \Gamma, U \rangle \circ \text{FrSt}; x.f := y \rightsquigarrow \mathbf{h}'; \Gamma^r}$	[OS-FIELD-WRITE]
$\begin{aligned} \iota &= \Gamma^r(y) \\ \iota &\neq \text{null}^r \\ C &= \text{class}^r(\iota) \\ \text{mLoc}(C, mt) &= \bar{v} \\ \text{mType}(C, mt) &= T_p \rightarrow T_r \\ \Gamma; U \vdash_{DE} y.mt(z) : U' \\ \Gamma^r_1 &= \{\langle p, \Gamma^r(z) \rangle, \overline{\langle v, \text{null}^r \rangle}, \langle \text{this}, \iota \rangle, \langle \text{res}, \text{null}^r \rangle\} \\ \mathbf{Fr}_1 &= \langle \Gamma^r_1, \Gamma_{C, mt}, \emptyset \rangle \\ \mathbf{Fr}_2 &= \langle \Gamma^r, \Gamma, U' \rangle \\ \mathbf{h}; \mathbf{Fr}_1 \circ \mathbf{Fr}_2 \circ \text{FrSt}; \mathbf{mBody}(C, mt) &\rightsquigarrow \mathbf{h}'; \Gamma^r_2 \\ \Gamma^r_3 &= \Gamma^r_1[x \mapsto \Gamma^r_2(\text{res})] \end{aligned}$	[OS-INVK]
$\frac{\mathbf{ow} = \text{mod20w}(\mathbf{h}, \text{mod}(T), \Gamma^r(\text{this})) \quad \langle \iota, \mathbf{h}' \rangle = \text{new}(\mathbf{h}, \text{class}(T), \mathbf{ow}) \quad \Gamma^r' = \Gamma^r[x \mapsto \iota]}{\mathbf{h}; \langle \Gamma^r, \Gamma, U \rangle \circ \text{FrSt}; x = \text{new } T() \rightsquigarrow \mathbf{h}'; \Gamma^r'}$	[OS-NEW]
$\frac{\mathbf{ow} = \text{mod20w}(\mathbf{h}, \text{mod}(T), \Gamma^r(\text{this})) \quad \text{mod}(T) \neq \text{any} \Rightarrow \mathbf{ow} = \text{owner}(\mathbf{h}, \Gamma^r(y)) \quad \text{class}^r(\Gamma^r(y)) \sqsubseteq \text{class}(T) \quad \Gamma^r' = \Gamma^r[x \mapsto \Gamma^r(y)]}{\mathbf{h}; \langle \Gamma^r, \Gamma, U \rangle \circ \text{FrSt}; x = (T)y \rightsquigarrow \mathbf{h}; \Gamma^r'}$	[OS-CAST]

$$\frac{\begin{array}{c} \mathbf{h}; \langle \Gamma^r, \Gamma, U \rangle \circ \mathbf{FrSt}; e_1 \rightsquigarrow \mathbf{h}_1; \Gamma^r_1 \\ \Gamma; U \vdash_{DE} e_1 : U' \end{array}}{\mathbf{h}; \langle \Gamma^r_1, \Gamma, U' \rangle \circ \mathbf{FrSt}; e_2 \rightsquigarrow \mathbf{h}_2; \Gamma^r_2} \quad [\text{OS-COMP}]$$

$$\frac{\begin{array}{c} \mathbf{ow} = \mathbf{mod20w}(\mathbf{h}, \mathbf{uniq}, \Gamma^r(\mathbf{this})) \\ \mathbf{h}' = \mathbf{transferCl}(\mathbf{h}, \Gamma^r(y), \mathbf{ow}) \\ \Gamma^{r'} = \Gamma^r[x \mapsto \Gamma^r(y)] \end{array}}{\mathbf{h}; \langle \Gamma^r, \Gamma, U \rangle \circ \mathbf{FrSt}; x = \mathbf{release}(y) \rightsquigarrow \mathbf{h}'; \Gamma^{r'}} \quad [\text{OS-RELEASE}]$$

$$\frac{\begin{array}{c} \mathbf{ow} = \mathbf{mod20w}(\mathbf{h}, m, \Gamma^r(\mathbf{this})) \\ \mathbf{h}' = \mathbf{transferCl}(\mathbf{h}, \Gamma^r(y), \mathbf{ow}) \\ \Gamma^{r'} = \Gamma^r[x \mapsto \Gamma^r(y)] \end{array}}{\mathbf{h}; \langle \Gamma^r, \Gamma, U \rangle \circ \mathbf{FrSt}; x = \mathbf{capture}\langle m \rangle(y) \rightsquigarrow \mathbf{h}'; \Gamma^{r'}} \quad [\text{OS-CAPTURE}]$$

6 Type safety

$\frac{}{\mathbf{h}; \iota_0 \vdash_{wf} \mathbf{null^r} : m}$	[WF-NULL]
$\frac{}{\mathbf{h}; \iota_0 \vdash_{wf} \iota : \mathbf{any}}$	[WF-ANY]
$\frac{\mathbf{owner}(\mathbf{h}, \iota) = \mathbf{owner}(\mathbf{h}, \iota_0)}{\mathbf{h}; \iota_0 \vdash_{wf} \iota : \mathbf{peer}}$	[WF-PEER]
$\frac{\iota = \iota_0}{\mathbf{h}; \iota_0 \vdash_{wf} \iota : \mathbf{this}}$	[WF-THIS]
$\frac{\mathbf{owner}(\mathbf{h}, \iota) = \langle Cl, \iota_0 \rangle}{\mathbf{h}; \iota_0 \vdash_{wf} \iota : \mathbf{rep} \langle Cl \rangle}$	[WF-REP]
$\frac{\mathbf{owner}(\mathbf{h}, \iota) = \langle Cl, \mathbf{null^r} \rangle \quad Cl \in \mathbf{CltFree}}{\mathbf{h}; \iota_0 \vdash_{wf} \iota : \mathbf{uniq}}$	[WF-UNIQ]
$\frac{\iota \neq \mathbf{null^r} \Rightarrow C \sqsubseteq \mathbf{class^r}(\mathbf{h}, \iota)}{\mathbf{h} \vdash_{wf} \iota : C}$	[WF-CLASS]
$\frac{\begin{array}{c} \mathbf{h} \vdash_{wf} \iota : C \\ \mathbf{h}; \iota_0 \vdash_{wf} \iota : m \end{array}}{\mathbf{h}; \iota_0 \vdash_{wf} \iota : m \ C}$	[WF-TYPE]
$\frac{\forall \iota \in \mathbf{dom}(\mathbf{h}) : C = \mathbf{class^r}(\mathbf{h}, \iota) \forall f \in \mathbf{fields}(C) : \mathbf{h} \vdash_{wf} \mathbf{h}(\iota.f) : \mathbf{class}(\mathbf{fType}(C, f))}{\vdash_{wf} \mathbf{h}}$	[WF-HEAP-CL]
$\frac{\forall \iota \in \mathbf{dom}(\mathbf{h}) : C = \mathbf{class^r}(\mathbf{h}, \iota) \forall f \in \mathbf{fields}(C) : \langle \iota, f \rangle \notin \{\iota_0\} \times U \Rightarrow \mathbf{h}; \iota \vdash_{wf} \mathbf{h}(\iota.f) : \mathbf{mod}(\mathbf{fType}(C, f))}{U; \iota_0 \vdash_{wf} \mathbf{h}}$	[WF-HEAP]

$\frac{\begin{array}{c} \text{dom}(\Gamma^r) = \text{dom}(\Gamma) \\ \text{dom}(\Gamma) \supseteq U \\ \text{mod}(\Gamma(\text{this})) = \text{this} \\ \forall v \in \text{dom}(\Gamma^r) \setminus U : h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^r(v) : \Gamma(v) \end{array}}{h \vdash_{wf} \Gamma^r; \Gamma; U}$	[WF-FRAME]
$\frac{\forall \iota \in \text{dom}(h) : \iota \notin \text{owners}(h, \iota) \downarrow_2}{\text{WfTree}(h)}$	[WF-TREE]
$\frac{\overline{h \vdash_{wf} \Gamma^r; \Gamma; U}}{h \vdash_{wf} \langle \Gamma^r, \Gamma, U \rangle}$	[WF-FRST]
$\frac{\forall Cl \in \text{CltFree} \text{ RefNum}(h, \text{FrSt}, \langle Cl, \text{null}^r \rangle, \text{uniq}) \leq 1}{h \vdash_{Uniq} \text{FrSt}}$	[WF-UNIQ-GLOB]
$\frac{\forall Cl \in \text{CltRep} \setminus \{Cl_{\text{this}}\} \text{ RefNum}(h, \text{FrSt}, \langle Cl, \iota \rangle, \text{rep } \langle Cl \rangle) = 0}{h; \text{FrSt} \vdash_{Inv} \iota}$	[WF-INV-OBJ]
$\frac{\begin{array}{c} (\{\text{owner}(h, \Gamma^r(\text{this}))\} \cup \\ \cup \{\text{owner}(h, \Gamma^r(x)) \mid \Gamma^r(x) \neq \text{null}^r \wedge \Gamma(x) = \text{uniq} \wedge x \notin U\}) \cap \\ \cap \text{owners}(h, \iota) \neq \emptyset \end{array}}{\text{ExtOwn}(h, \langle \Gamma^r, \Gamma, U \rangle, \iota)}$	[EXT-OWN-TR]
$\frac{\begin{array}{c} \text{FrSt} = \langle \Gamma^r, \Gamma, U \rangle \circ \text{FrSt}' \\ \iota_0 = \Gamma^r(\text{this}) \\ h; \text{FrSt}' \vdash_{Inv} \iota_0 \\ \forall \iota : \text{ExtOwn}(h, \text{Fr}, \iota) \wedge \iota \neq \iota_0 \Rightarrow h; \text{FrSt} \vdash_{Inv} \iota \end{array}}{h \vdash_{Inv} \text{FrSt}}$	[WF-INV-GLOB]
$\frac{\begin{array}{c} \text{Fr}_0 \circ \overline{\text{Fr}} = \langle \Gamma_0^r, \Gamma_0, U_0 \rangle \circ \overline{\langle \Gamma^r, \Gamma, U \rangle} \\ \forall Cl \in \text{CltFree} : \langle Cl, \text{null}^r \rangle \notin \text{owners}(h, \Gamma_0^r(\text{this})) \\ \text{owner}(h, \Gamma^r(\text{this})) \in \text{owners}(h, \Gamma_0^r(\text{this})) \end{array}}{\text{ThisOw}(h, \text{Fr}_0 \circ \overline{\text{Fr}})}$	[WF-THIS-OW]

$$\begin{array}{c}
\text{FrSt} = \langle \Gamma^r, \Gamma, U \rangle \circ \text{FrSt}' \\
\vdash_{wf} h \\
U; \Gamma^r(\text{this}) \vdash_{wf} h \\
h \vdash_{wf} \text{FrSt} \\
h \vdash_{Uniq} \text{FrSt} \\
h \vdash_{Inv} \text{FrSt} \\
\text{WfTree}(h) \\
\text{ThisOwn}(h, \text{FrSt}) \\
\hline
\vdash_{wf} \langle h, \text{FrSt} \rangle
\end{array} \quad [\text{WF-ST}]$$

$$\frac{\forall \iota \in \text{dom}(h) : \neg \text{ExtOwn}(h, \text{Fr}, \iota) \Rightarrow h(\iota) = h'(\iota) \wedge \neg \text{ExtOwn}(h', \text{Fr}', \iota)}{\text{OwAsMod}(h, h', \text{Fr}, \text{Fr}')} \quad [\text{Ow-AS-MOD}]$$

7 Type safety proof

7.1 The main theorem

Theorem 7.1.

$$\left. \begin{array}{l}
\vdash_{wf} \langle h, \text{FrSt} \rangle \\
\text{FrSt} = \text{Fr} \circ \text{FrSt}_0 \\
\text{Fr} = \langle \Gamma^r, \Gamma, U \rangle \\
\Gamma; U \vdash e \\
\Gamma; U \vdash_{DE} e : U' \\
h; \text{FrSt}; e \rightsquigarrow h'; \Gamma^{r'} \\
\text{Fr}' = \langle \Gamma^{r'}, \Gamma, U' \rangle \\
\text{FrSt}' = \text{Fr}' \circ \text{FrSt}_0
\end{array} \right\} \Rightarrow \left\{ \begin{array}{l}
\vdash_{wf} \langle h', \text{FrSt}' \rangle \\
\text{OwAsMod}(h, h', \text{Fr}, \text{Fr}'')
\end{array} \right.$$

Proof.

$$\vdash_{wf} \langle h, \text{FrSt} \rangle \Rightarrow \left\{ \begin{array}{ll}
\vdash_{wf} h & (1) \\
U; \Gamma^r(\text{this}) \vdash_{wf} h & (2) \\
h \vdash_{wf} \text{FrSt} & (3) \\
h \vdash_{Uniq} \text{FrSt} & (4) \\
h \vdash_{Inv} \text{FrSt} & (5) \\
\text{WfTree}(h) & (6) \\
\text{ThisOwn}(h, \text{FrSt}) & (7)
\end{array} \right.$$

- OS-ASSIGN

$$U' = \text{consumeUniq}(U, \Gamma(x), y) \setminus \{x\}$$

$$\Gamma^{r'} = \Gamma^r[x \mapsto \Gamma^r(y)]$$

$$(2) \xrightarrow{\text{Lem. 7.14}} U'; \Gamma^{r'}(\text{this}) \vdash_{wf} h \quad (8)$$

$$\text{T-ASSIGN} \Rightarrow y \notin U \xrightarrow{(3)} h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^r(y) : \Gamma(y) \Rightarrow$$

$$\xrightarrow{\text{Lem. 7.21} \wedge \text{Lem. 7.22} \wedge (3)} h \vdash_{wf} \text{FrSt}' \quad (9)$$

$$\text{FrSt}'' = \langle \Gamma^r, \Gamma, \text{consumeUniq}(U, \Gamma(x), y) \rangle \circ \text{FrSt}_0$$

$$(4) \xrightarrow{\text{Lem. 7.29} \wedge (3)} h \vdash_{Uniq} \text{FrSt}'' \quad (10)$$

$$(3) \wedge (4) \Rightarrow (\Gamma^r(y) \neq \text{null}^r \wedge \Gamma(x) = \text{uniq} \Rightarrow$$

$$\text{RefNum}(h, \text{FrSt}'', \text{owner}(h, (\Gamma^r(y)), \text{uniq}) = 0)) \xrightarrow{\text{Lem. 7.30} \wedge (10)} h \vdash_{Uniq} \text{FrSt}' \quad (11)$$

$$(5) \wedge (9) \xrightarrow{\text{Lem. 7.37}} h \vdash_{Inv} \text{FrSt}' \quad (12)$$

$$(7) \xrightarrow{\text{Lem. 7.49}} \text{ThisOw}(h, \text{FrSt}') \quad (13)$$

$$(1) \wedge (8) \wedge (9) \wedge (11) \wedge (12) \wedge (6) \wedge (13) \Rightarrow \vdash_{wf} \langle h', \text{FrSt}' \rangle$$

Lem. 7.55 $\Rightarrow \text{OwAsMod}(h, h, \text{Fr}, \text{Fr}')$

- OS-FIELD-READ

$$U' = U \setminus \{x\}$$

$$\Gamma^r(y) \neq \text{null}^r$$

$$\Gamma^{r'} = \Gamma^r[x \mapsto h(\Gamma^r(y) . f)]$$

$$(2) \xrightarrow{\text{Lem. 7.14}} U'; \Gamma^{r'}(\text{this}) \vdash_{wf} h \quad (8)$$

$$\text{T-FIELD-READ} \Rightarrow y \notin U \xrightarrow{(3)} h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^r(y) : \Gamma(y) \xrightarrow{\text{Lem. 7.15}} h; \Gamma^r(\text{this}) \vdash_{wf} h(\Gamma^r(y) . f) : T \triangleright_U \text{fType}(\text{class}^r(\Gamma^r(y)), f) \Rightarrow$$

$$\xrightarrow{\text{Lem. 7.21} \wedge \text{Lem. 7.22} \wedge (3)} h \vdash_{wf} \text{FrSt}' \quad (9)$$

$$\text{T-FIELD-READ} \Rightarrow \Gamma(x) \neq \text{uniq} \xrightarrow{\text{Lem. 7.30} \wedge \text{Lem. 7.29} \wedge (4)} h \vdash_{Uniq} \text{FrSt}' \quad (10)$$

$$(5) \wedge (9) \xrightarrow{\text{Lem. 7.37}} h \vdash_{Inv} \text{FrSt}' \quad (11)$$

$$(7) \xrightarrow{\text{Lem. 7.49}} \text{ThisOw}(h, \text{FrSt}') \quad (12)$$

$$(1) \wedge (8) \wedge (9) \wedge (10) \wedge (11) \wedge (6) \wedge (12) \Rightarrow \vdash_{wf} \langle h', \text{FrSt}' \rangle$$

Lem. 7.55 $\Rightarrow \text{OwAsMod}(h, h, \text{Fr}, \text{Fr}')$

- OS-FIELD-WRITE

$U' = \text{if}(y = \text{this}) \text{ then } U \setminus \{f\} \text{ else } U$
 $\mathbf{h}' = \mathbf{h}[\Gamma^{\mathbf{r}}(x) . f := \Gamma^{\mathbf{r}}(y)]$

- (1) $\xrightarrow{\text{Lem. 7.11}} \vdash_{wf} \mathbf{h}'$ (8)
- (3) $\Rightarrow \mathbf{h}; \Gamma^{\mathbf{r}}(\text{this}) \vdash_{wf} \Gamma^{\mathbf{r}}(y) : \Gamma(y) \wedge \mathbf{h}; \Gamma^{\mathbf{r}}(\text{this}) \vdash_{wf} \Gamma^{\mathbf{r}}(x) : \Gamma(x) \Rightarrow$
 $\xrightarrow{\text{Lem. 7.3} \wedge \text{T-FIELD-WRITE}} \mathbf{h}; \Gamma^{\mathbf{r}}(y) \vdash_{wf} \Gamma^{\mathbf{r}}(x) : \text{fType}(\text{class}^{\mathbf{r}}(\Gamma^{\mathbf{r}}(x)), f) \Rightarrow$
 $\xrightarrow{\text{Lem. 7.15} \wedge (2)} U'; \Gamma^{\mathbf{r}'}(\text{this}) \vdash_{wf} \mathbf{h}'$ (9)
- (3) $\xrightarrow{\text{Lem. 7.21} \wedge \text{Lem. 7.23}} \mathbf{h}' \vdash_{wf} \text{FrSt}'$ (10)
- (4) $\xrightarrow{\text{Lem. 7.29} \wedge \text{Lem. 7.31}} \mathbf{h}' \vdash_{Uniq} \text{FrSt}'$ (11)
- (5) $\xrightarrow{\text{Lem. 7.37} \wedge \text{Lem. 7.38}} \mathbf{h}' \vdash_{Inv} \text{FrSt}'$ (12)
- (6) $\xrightarrow{\text{Lem. 7.45}} \text{WfTree}(\mathbf{h}')$ (13)
- (7) $\xrightarrow{\text{Lem. 7.49} \wedge \text{Lem. 7.50}} \text{This} \mathbf{0w}(\mathbf{h}', \text{FrSt}')$ (14)
- (8) \wedge (9) \wedge (10) \wedge (11) \wedge (12) \wedge (13) \wedge (14) $\Rightarrow \vdash_{wf} \langle \mathbf{h}', \text{FrSt}' \rangle$
 $\text{Lem. 7.55} \wedge \text{Lem. 7.56} \Rightarrow \mathbf{0wAsMod}(\mathbf{h}, \mathbf{hFr}', \mathbf{Fr}, \mathbf{Fr}')$

- OS-INVK

$$\begin{aligned}
U_1 &= \text{consumeUniq}(U, T_p, z) \\
U_2 &= \text{consumeLocals}(U_1) \\
\Gamma^r_1 &= \{\langle p, \Gamma^r(z) \rangle, \overline{\langle v, \text{null}^r \rangle}, \langle \text{this}, \iota \rangle, \langle \text{res}, \text{null}^r \rangle\} \\
\mathbf{Fr}_1 &= \langle \Gamma^r_1, \Gamma_C, mt, \emptyset \rangle \\
\mathbf{Fr}_2 &= \langle \Gamma^r, \Gamma, U_2 \rangle \\
\mathbf{FrSt}_1 &= \mathbf{Fr}_1 \circ \mathbf{Fr}_2 \circ \mathbf{FrSt}_0 \\
h; \mathbf{FrSt}_1; \mathbf{mBody}(C, mt) &\rightsquigarrow h'; \Gamma^r_2 \\
\Gamma; U_2 \vdash_{DE} y.mt(z) : U_3 \\
\mathbf{Fr}_3 &= \langle \Gamma^r_2, \Gamma_C, mt, U_3 \rangle \\
\mathbf{FrSt}_2 &= \mathbf{Fr}_3 \circ \mathbf{Fr}_2 \circ \mathbf{FrSt}_0 \\
\Gamma^r_3 &= \Gamma^r_1[x \mapsto \Gamma^r_2(\text{res})] \\
U_4 &= U_2 \setminus \{x\} \\
\mathbf{Fr}' &= \langle \Gamma^r_3, \Gamma_C, mt, U_4 \rangle
\end{aligned}$$

$$\begin{aligned}
(2) &\stackrel{\text{Lem. 7.14}}{\Rightarrow} U_2; \Gamma^r(\text{this}) \vdash_{wf} h \stackrel{\text{Lem. 7.17}}{\Rightarrow} \emptyset; \Gamma^r_1(\text{this}) \vdash_{wf} h \quad (9) \\
(3) \wedge \text{T-INVK} &\stackrel{\text{Lem. 7.21} \wedge \text{Lem. 7.24}}{\Rightarrow} h \vdash_{wf} \mathbf{FrSt}_1 \quad (10) \\
(4) &\stackrel{\text{Lem. 7.29} \wedge \text{Lem. 7.32}}{\Rightarrow} (5) \stackrel{\text{Lem. 7.37} \wedge \text{Lem. 7.39}}{\Rightarrow} h \vdash_{Inv} \mathbf{FrSt}_1 \quad (12) \\
(7) &\stackrel{\text{Lem. 7.49} \wedge \text{Lem. 7.51}}{\Rightarrow} \text{ThisOw}(h, \mathbf{FrSt}_1) \quad (13) \\
(1) \wedge (9) \wedge (10) \wedge (11) \wedge (12) \wedge (6) \wedge (13) &\Rightarrow \vdash_{wf} \langle h, \mathbf{FrSt}_2 \rangle \stackrel{\text{Ind. step}}{\Rightarrow} \\
&\Rightarrow \left\{ \begin{array}{ll} \vdash_{wf} \langle h', \mathbf{FrSt}_2 \rangle & (14) \\ \text{OwAsMod}(h, h', \mathbf{Fr}_1, \mathbf{Fr}_3) & (15) \end{array} \right. \\
(14) \Rightarrow &\left\{ \begin{array}{ll} \vdash_{wf} h' & (16) \\ \mathbf{FrSt}_2; h' \vdash_{wf} & (17) \\ U_3; \Gamma^r_2(\text{this}) \vdash_{wf} h & (18) \\ h' \vdash_{Uniq} \mathbf{FrSt}_2 & (19) \\ h' \vdash_{Inv} \mathbf{FrSt}_2 & (20) \\ \text{WfTree}(h') & (21) \\ \text{ThisOw}(h', \mathbf{FrSt}_2) & (22) \end{array} \right.
\end{aligned}$$

$$\begin{aligned}
T\text{-INVK} \Rightarrow U_3 \cap \text{FieldID} = \emptyset &\stackrel{\text{Lem. 7.17}\wedge(17)}{\Rightarrow} U_2; \Gamma^r(\text{this}) \vdash_{wf} h' \Rightarrow \\
&\stackrel{\text{Lem. 7.14}}{\Rightarrow} U_4; \Gamma^r_3(\text{this}) \vdash_{wf} h' \quad (23) \\
(18) &\stackrel{\text{Lem. 7.25}}{\Rightarrow} h' \vdash_{wf} \text{Fr}_2 \circ \text{FrSt}_0 \stackrel{\text{Lem. 7.22}}{\Rightarrow} h' \vdash_{wf} \text{FrSt}' \quad (24) \\
(19) &\stackrel{\text{Lem. 7.33}}{\Rightarrow} h' \vdash_{\text{Uniq}} \text{Fr}_2 \circ \text{FrSt}_0 \stackrel{\text{Lem. 7.30}}{\Rightarrow} h' \vdash_{\text{Uniq}} \text{FrSt}' \quad (25) \\
(20) \wedge (15) &\stackrel{\text{Lem. 7.40}\wedge\text{Lem. 7.41}}{\Rightarrow} h' \vdash_{\text{Inv}} \text{Fr}_2 \circ \text{FrSt}_0 \stackrel{\text{Lem. 7.37}}{\Rightarrow} h' \vdash_{\text{Inv}} \text{FrSt}' \quad (26) \\
(22) \wedge (15) &\stackrel{\text{Lem. 7.53}\wedge\text{Lem. 7.53}}{\Rightarrow} \text{This}0w(h', \text{Fr}_2 \circ \text{FrSt}_0) \stackrel{\text{Lem. 7.49}}{\Rightarrow} \text{This}0w(h', \text{FrSt}') \quad (27) \\
(16) \wedge (23) \wedge (24) \wedge (25) \wedge (26) \wedge (21) \wedge (27) &\Rightarrow \vdash_{wf} \langle h', \text{FrSt}' \rangle \\
\text{Lem. 7.55} &\Rightarrow \text{OwAsMod}(h, h, \text{Fr}, \text{Fr}_2) \quad (28) \\
\text{Lem. 7.55} &\Rightarrow \text{OwAsMod}(h, h, \text{Fr}_2, \text{Fr}_1) \stackrel{\text{Lem. 7.59}\wedge(15)}{\Rightarrow} \text{OwAsMod}(h, h', \text{Fr}_2, \text{Fr}_3) \Rightarrow \\
&\stackrel{\text{Lem. 7.57}}{\Rightarrow} \text{OwAsMod}(h, h', \text{Fr}_2, \text{Fr}_2) \quad (29) \\
(28) \wedge (29) &\stackrel{\text{Lem. 7.59}}{\Rightarrow} \text{OwAsMod}(h, h', \text{Fr}, \text{Fr}_2) \stackrel{\text{Lem. 7.55}}{\Rightarrow} \text{OwAsMod}(h, h', \text{Fr}, \text{Fr}')
\end{aligned}$$

- OS-NEW

$$\begin{aligned}
U' &= U \setminus \{x\} \\
\text{ow} &= \text{mod20w}(h, \text{mod}(T), \Gamma^r(\text{this})) \\
\langle \iota, h' \rangle &= \text{new}(h, \text{class}(T), \text{ow}) \\
\Gamma^{r'} &= \Gamma^r[x \mapsto \iota]
\end{aligned}$$

$$\begin{aligned}
(1) &\stackrel{\text{Lem. 7.12}}{\Rightarrow} \vdash_{wf} h' \quad (8) \\
(2) &\stackrel{\text{Lem. 7.14}}{\Rightarrow} U'; \Gamma^{r'}(\text{this}) \vdash_{wf} h \stackrel{\text{Lem. 7.18}}{\Rightarrow} U'; \Gamma^{r'}(\text{this}) \vdash_{wf} h' \quad (9) \\
(3) &\stackrel{\text{Lem. 7.26}\wedge(3)}{\Rightarrow} h' \vdash_{wf} \text{FrSt} \quad (10) \\
\text{Lem. 7.4} \Rightarrow h; \Gamma^r(\text{this}) \vdash_{wf} \iota : T &\stackrel{\text{Lem. 7.6}\wedge\text{T-NEW}}{\Rightarrow} h; \Gamma^r(\text{this}) \vdash_{wf} \iota : \Gamma(x) \Rightarrow \\
&\stackrel{\text{Lem. 7.22}\wedge(10)}{\Rightarrow} h' \vdash_{wf} \text{FrSt}' \quad (11) \\
\text{mod}(T) = \text{uniq} \Rightarrow \text{ow} &= \langle \text{newClt}(h), \text{null}^r \rangle \Rightarrow \text{RefNum}(h', \text{FrSt}, \text{ow}, \text{uniq}) = 0 \quad (12) \\
(4) &\stackrel{\text{Lem. 7.34}}{\Rightarrow} h' \vdash_{\text{Uniq}} \text{FrSt} \stackrel{\text{Lem. 7.30}\wedge(12)}{\Rightarrow} h' \vdash_{\text{Uniq}} \text{FrSt}' \quad (13) \\
(5) &\stackrel{\text{Lem. 7.42}}{\Rightarrow} h' \vdash_{\text{Inv}} \text{FrSt} \stackrel{\text{Lem. 7.37}}{\Rightarrow} h' \vdash_{\text{Inv}} \text{FrSt}' \quad (13) \\
(6) &\stackrel{\text{Lem. 7.46}}{\Rightarrow} \text{WfTree}(h') \quad (14) \\
(7) &\stackrel{\text{Lem. 7.52}}{\Rightarrow} \text{This}0w(h', \text{FrSt}) \stackrel{\text{Lem. 7.49}}{\Rightarrow} \text{This}0w(h', \text{FrSt}') \quad (15) \\
(8) \wedge (9) \wedge (10) \wedge (12) \wedge (13) \wedge (14) \wedge (15) &\Rightarrow \vdash_{wf} \langle h', \text{FrSt}' \rangle \\
\text{Lem. 7.58} &\Rightarrow \text{OwAsMod}(h, h', \text{FrSt}, \text{FrSt}')
\end{aligned}$$

- OS-CAST

$$\begin{aligned}
U' &= \text{consumeUniq}(U, \Gamma(x), y) \setminus \{x\} \\
\text{ow} &= \text{mod20w}(\text{h}, \text{mod}(T), \Gamma^{\text{r}}(\text{this})) \\
\text{mod}(T) \neq \text{any} &\Rightarrow \text{ow} = \text{owner}(\text{h}, \Gamma^{\text{r}}(y)) \\
\text{class}^{\text{r}}(\Gamma^{\text{r}}(y)) &\sqsubseteq \text{class}(T) \\
\Gamma^{\text{r}'} &= \Gamma^{\text{r}}[x \mapsto \Gamma^{\text{r}}(y)]
\end{aligned}$$

$$\begin{aligned}
(2) &\stackrel{\text{Lem. 7.14}}{\Rightarrow} U'; \Gamma^{\text{r}'}(\text{this}) \vdash_{wf} \text{h} \quad (8) \\
(3) &\stackrel{\text{Lem. 7.4}}{\Rightarrow} \text{h}; \Gamma^{\text{r}}(\text{this}) \vdash_{wf} \Gamma^{\text{r}}(y) : T \Rightarrow \\
&\stackrel{\text{Lem. 7.21} \wedge \text{Lem. 7.22} \wedge (3)}{\Rightarrow} \text{h} \vdash_{wf} \text{FrSt}' \quad (9) \\
\text{FrSt}'' &= \langle \Gamma^{\text{r}}, \Gamma, \text{consumeUniq}(U, \Gamma(x), y) \rangle \circ \text{FrSt}_0 \\
(4) &\stackrel{\text{Lem. 7.29} \wedge (3)}{\Rightarrow} \text{h} \vdash_{Uniq} \text{FrSt}'' \quad (10) \\
(3) \wedge (4) &\Rightarrow (\Gamma^{\text{r}}(y) \neq \text{null}^{\text{r}} \wedge \Gamma(x) = \text{uniq} \Rightarrow \\
\text{RefNum}(\text{h}, \text{FrSt}'', \text{owner}(\text{h}, (\Gamma^{\text{r}}(y)), \text{uniq}) = 0)) &\stackrel{\text{Lem. 7.30} \wedge (10)}{\Rightarrow} \text{h} \vdash_{Uniq} \text{FrSt}' \quad (11) \\
(5) \wedge (9) &\stackrel{\text{Lem. 7.37}}{\Rightarrow} \text{h} \vdash_{Inv} \text{FrSt}' \quad (12) \\
(7) &\stackrel{\text{Lem. 7.49}}{\Rightarrow} \text{This} \text{ow}(\text{h}, \text{FrSt}') \quad (13) \\
(1) \wedge (8) \wedge (9) \wedge (11) \wedge (12) \wedge (6) \wedge (13) &\Rightarrow \vdash_{wf} \langle \text{h}', \text{FrSt}' \rangle \\
\text{Lem. 7.55} &\Rightarrow \text{OwAsMod}(\text{h}, \text{h}, \text{Fr}, \text{Fr})
\end{aligned}$$

- OS-COMP

$$\begin{aligned}
\Gamma; U &\vdash_{DE} e_1; e_2 : U'' \\
\text{Fr}_1 &= \langle \Gamma^{\text{r}}_1, \Gamma, U' \rangle \\
\vdash_{wf} \langle \text{h}, \text{FrSt} \rangle &\stackrel{\text{Ind.}}{\Rightarrow} \stackrel{\text{step}}{\left\{ \begin{array}{l} \vdash_{wf} \langle \text{h}_1, \text{Fr}_1 \circ \text{FrSt} \rangle \\ \text{OwAsMod}(\text{h}, \text{h}_1, \text{Fr}, \text{Fr}_1) \end{array} \right\}} \quad (7) \\
&\quad (8) \\
\text{Fr}_2 &= \langle \Gamma^{\text{r}}_2, \Gamma, U'' \rangle \\
(7) &\stackrel{\text{Ind.}}{\Rightarrow} \stackrel{\text{step}}{\left\{ \begin{array}{l} \vdash_{wf} \langle \text{h}_2, \text{Fr}_2 \circ \text{FrSt} \rangle \\ \text{OwAsMod}(\text{h}_1, \text{h}_2, \text{Fr}_1, \text{Fr}_2) \end{array} \right\}} \quad (9) \\
(8) \wedge (9) &\stackrel{\text{Lem. 7.59}}{\Rightarrow} \text{OwAsMod}(\text{h}, \text{h}_2, \text{Fr}_1, \text{Fr}_2)
\end{aligned}$$

- OS-RELEASE

$$\begin{aligned}
U' &= \text{consumeAliases}(U, y) \setminus \{x\} \\
\text{ow} &= \text{mod20w}(h, \text{uniq}, \Gamma^r(\text{this})) \\
h' &= \text{transferCl}(h, \Gamma^r(y), \text{ow}) \\
\Gamma^{r'} &= \Gamma^r[x \mapsto \Gamma^r(y)]
\end{aligned}$$

- (1) $\stackrel{\text{Lem. 7.13}}{\Rightarrow} \vdash_{wf} h' \quad (8)$
- (2) $\stackrel{\text{Lem. 7.19} \wedge \text{Lem. 7.14}}{\Rightarrow} U'; \Gamma^{r'}(\text{this}) \vdash_{wf} h' \quad (9)$
- (3) \wedge (5) $\stackrel{\text{Lem. 7.27} \wedge \text{Lem. 7.21}}{\Rightarrow} h' \vdash_{wf} \text{FrSt}' \quad (10)$
- (3) \wedge (4) $\stackrel{\text{Lem. 7.35} \wedge \text{Lem. 7.29}}{\Rightarrow} h' \vdash_{\text{Uniq}} \text{FrSt}' \quad (11)$
- (3) \wedge (5) \wedge (7) $\stackrel{\text{Lem. 7.43} \wedge \text{Lem. 7.37}}{\Rightarrow} h' \vdash_{\text{Inv}} \text{FrSt}' \quad (12)$
- (6) $\stackrel{\text{Lem. 7.47}}{\Rightarrow} \text{WfTree}(h') \quad (13)$
- (7) $\stackrel{\text{Lem. 7.54} \wedge \text{Lem. 7.49}}{\Rightarrow} \text{This0w}(h', \text{FrSt}') \quad (14)$
- (8) \wedge (9) \wedge (10) \wedge (11) \wedge (12) \wedge (13) \wedge (14) $\Rightarrow \vdash_{wf} \langle h', \text{FrSt}' \rangle$
- (3) \wedge Lem. 7.60 \wedge Lem. 7.55 $\Rightarrow \text{0wAsMod}(h, h', \text{Fr}, \text{Fr}')$

• OS-CAPTURE

$$\begin{aligned}
U' &= U \cup \{y\} \setminus \{x\} \\
\text{ow} &= \text{mod20w}(h, m, \Gamma^r(\text{this})) \\
h' &= \text{transferCl}(h, \Gamma^r(y), \text{ow}) \\
\Gamma^{r'} &= \Gamma^r[x \mapsto \Gamma^r(y)]
\end{aligned}$$

- (1) $\stackrel{\text{Lem. 7.13}}{\Rightarrow} \vdash_{wf} h' \quad (8)$
- (2) $\stackrel{\text{Lem. 7.20} \wedge \text{Lem. 7.14}}{\Rightarrow} U'; \Gamma^{r'}(\text{this}) \vdash_{wf} h' \quad (9)$
- (3) \wedge (5) $\stackrel{\text{Lem. 7.28} \wedge \text{Lem. 7.21}}{\Rightarrow} h' \vdash_{wf} \text{FrSt}' \quad (10)$
- (3) \wedge (4) $\stackrel{\text{Lem. 7.36} \wedge \text{Lem. 7.29}}{\Rightarrow} h' \vdash_{\text{Uniq}} \text{FrSt}' \quad (11)$
- (3) \wedge (5) \wedge (7) $\stackrel{\text{Lem. 7.44} \wedge \text{Lem. 7.37}}{\Rightarrow} h' \vdash_{\text{Inv}} \text{FrSt}' \quad (12)$
- (6) $\stackrel{\text{Lem. 7.48}}{\Rightarrow} \text{WfTree}(h') \quad (13)$
- (7) $\stackrel{\text{Lem. 7.54} \wedge \text{Lem. 7.49}}{\Rightarrow} \text{This0w}(h', \text{FrSt}') \quad (14)$
- (8) \wedge (9) \wedge (10) \wedge (11) \wedge (12) \wedge (13) \wedge (14) $\Rightarrow \vdash_{wf} \langle h', \text{FrSt}' \rangle$
- (3) \wedge Lem. 7.61 \wedge Lem. 7.55 $\Rightarrow \text{0wAsMod}(h, h', \text{Fr}, \text{Fr}')$

□

7.2 Auxiliary lemmas

$$\begin{aligned} \text{RepSet} &:: \text{Heap} \times \text{Addr!} \rightarrow \mathcal{P}(\text{Addr!}) \\ \text{RepSet}(h, \iota) &= \{\iota' \in \text{dom}(h) \mid \text{owner}(h, \iota') \downarrow_2 = \iota\} \end{aligned}$$

$$\begin{aligned} \text{RepTrSet} &:: \text{Heap} \times \text{Addr!} \rightarrow \mathcal{P}(\text{Addr!}) \\ \text{RepTrSet}(h, \iota) &= \{\iota' \in \text{dom}(h) \mid \iota \in \text{owners}(h, \iota') \downarrow_2\} \end{aligned}$$

$$\begin{aligned} \text{PeerSet} &:: \text{Heap} \times \text{Addr!} \rightarrow \mathcal{P}(\text{Addr!}) \\ \text{PeerSet}(h, \iota) &= \{\iota' \in \text{dom}(h) \mid \text{owner}(h, \iota) = \text{owner}(h, \iota')\} \end{aligned}$$

Lemma 7.1.

$$h; \iota_0 \vdash_{wf} \iota : \text{uniq} \Rightarrow h; \iota'_0 \vdash_{wf} \iota : \text{uniq}$$

Proof.

$$\begin{aligned} h; \iota_0 \vdash_{wf} \iota : \text{uniq} &\Rightarrow \text{owner}(h, \iota) = \langle Cl, \text{null}^r \rangle \wedge \\ Cl \in \text{CltFree} &\Rightarrow h; \iota'_0 \vdash_{wf} \iota : \text{uniq} \end{aligned}$$

□

Lemma 7.2.

$$\left. \begin{array}{l} h; \iota_0 \vdash_{wf} \iota : T \\ h; \iota \vdash_{wf} \iota' : T' \end{array} \right\} \stackrel{(1)}{\quad} \stackrel{(2)}{\quad} \Rightarrow h; \iota_0 \vdash_{wf} \iota' : T \triangleright_U T'$$

Proof.

$$\begin{aligned} T &= C \ m \\ T' &= C' \ m' \\ (2) \Rightarrow C' &\sqsubseteq \text{class}^r(h, \ i') \ (3) \end{aligned}$$

$$\begin{aligned} m = \text{peer} \wedge m' = \text{peer} &\xrightarrow{(1)\wedge(2)} \\ \text{owner}(h, \ i_0) = \text{owner}(h, \ i) \wedge \text{owner}(h, \ i) = \text{owner}(h, \ i') \Rightarrow \\ \text{owner}(h, \ i_0) = \text{owner}(h, \ i') \Rightarrow h; i_0 \vdash_{wf} i' : \text{peer} \ (4) \end{aligned}$$

$$\begin{aligned} m = \text{rep} \langle Cl \rangle \wedge m' = \text{peer} &\xrightarrow{(1)\wedge(2)} \\ \text{owner}(h, \ i) = \langle i_0, \ Cl \rangle \wedge \text{owner}(h, \ i) = \text{owner}(h, \ i') \Rightarrow \\ \langle i_0, \ Cl \rangle = \text{owner}(h, \ i') \Rightarrow h; i_0 \vdash_{wf} i' : \text{rep} \langle Cl \rangle \ (5) \end{aligned}$$

$$m = \text{this} \xrightarrow{(1)} i = i_0 \xrightarrow{(2)} h; i_0 \vdash_{wf} i' : m' \ (6)$$

$$m' = \text{uniq} \xrightarrow{(2)\wedge\text{Lem. 7.1}} h; i_0 \vdash_{wf} i' : \text{uniq} \ (7)$$

$$\text{In all other cases } m \triangleright_U m' = \text{any} \Rightarrow h; i_0 \vdash_{wf} i' : \text{any} \ (8)$$

$$(4) \wedge (5) \wedge (6) \wedge (7) \wedge (8) \Rightarrow h; i_0 \vdash_{wf} i' : m \triangleright_U m' \ (9)$$

$$(2) \wedge (9) \Rightarrow h; i_0 \vdash_{wf} i' : T \triangleright_U T'$$

□

Lemma 7.3.

$$\left. \begin{array}{l} h; i_0 \vdash_{wf} i : T \\ h; i_0 \vdash_{wf} i' : T \triangleright_U T' \\ \text{IsWritable}(T) \\ T' = \text{rep} \langle Cl \rangle \Rightarrow T = \text{this} \end{array} \right\} \begin{array}{l} (1) \\ (2) \\ (3) \\ (4) \end{array} \Rightarrow h; i \vdash_{wf} i' : T'$$

Proof.

$$\begin{aligned} T &= C \ m \\ T' &= C' \ m' \\ (2) \Rightarrow C' &\sqsubseteq \text{class}^r(h, \ i') \ (5) \end{aligned}$$

$$\begin{aligned} m = \text{peer} \wedge m' = \text{peer} &\xrightarrow{(1)\wedge(2)} \\ \text{owner}(h, \ i_0) = \text{owner}(h, \ i) \wedge \text{owner}(h, \ i_0) = \text{owner}(h, \ i') \Rightarrow \\ \text{owner}(h, \ i) = \text{owner}(h, \ i') \Rightarrow h; i \vdash_{wf} i' : \text{peer} \ (6) \end{aligned}$$

$$\begin{aligned} m = \text{rep} \langle Cl \rangle \wedge m' = \text{peer} &\xrightarrow{(1)\wedge(2)} \\ \text{owner}(h, \ i) = \langle i_0, \ Cl \rangle \wedge \text{owner}(h, \ i)' = \langle i_0, \ Cl \rangle \Rightarrow \\ \text{owner}(h, \ i) = \text{owner}(h, \ i') \Rightarrow h; i \vdash_{wf} i' : \text{peer} \ (7) \end{aligned}$$

$$m = \text{this} \xrightarrow{(1)} i = i_0 \xrightarrow{(2)} h; i \vdash_{wf} i' : m' \ (8)$$

$$m' = \text{uniq} \xrightarrow{(2)\wedge\text{Lem. 7.1}} h; i \vdash_{wf} i' : \text{uniq} \ (9)$$

$$m' = \text{any} \Rightarrow h; i \vdash_{wf} i' : \text{any} \ (10)$$

$m \setminus m'$	$\text{rep} \langle \cdot \rangle$	peer	any	uniq
$\text{rep} \langle \cdot \rangle$	(4)	(7)	(10)	(9)
peer	(4)	(6)	(10)	(9)
this	(8)	(8)	(8)(10)	(8)(9)
any	(3)(4)	(3)	(3)(10)	(3)(9)
uniq	(3)(4)	(3)	(3)(10)	(3)(9)

$$\Rightarrow h; i \vdash_{wf} i' : m' \ (11)$$

$$(5) \wedge (11) \Rightarrow h; i \vdash_{wf} i' : T'$$

□

Lemma 7.4.

$$\left. \begin{array}{l} \text{owner}(h, \ i) = \text{mod20w}(h, \ m, \ i_0) \quad (1) \\ m \neq \text{any} \quad (2) \end{array} \right\} \Rightarrow h; i_0 \vdash_{wf} i : m$$

Proof.

$$\begin{aligned}
 m = \text{rep} \langle Cl \rangle &\stackrel{(1)}{\Rightarrow} \text{owner}(h, \iota) = \langle Cl, \iota \rangle \Rightarrow h; \iota_0 \vdash_{wf} \iota : \text{rep} \langle Cl \rangle \quad (3) \\
 m = \text{peer} &\stackrel{(1)}{\Rightarrow} \text{owner}(h, \iota) = \text{owner}(h, \iota_0) \Rightarrow h; \iota_0 \vdash_{wf} \iota : \text{peer} \quad (4) \\
 m = \text{uniq} &\stackrel{(1)}{\Rightarrow} \langle Cl, \text{null}^r \rangle \wedge Cl \in \text{CltFree} \Rightarrow h; \iota_0 \vdash_{wf} \iota : \text{uniq} \quad (5) \\
 (2) \wedge (3) \wedge (4) \wedge (5) &\Rightarrow h; \iota_0 \vdash_{wf} \iota : m
 \end{aligned}$$

□

Lemma 7.5.

$$\left. \begin{array}{l} \text{class}^r(h, \iota) = C \\ h' = \text{transfer}(h, \iota_{tr}, \text{ow}) \end{array} \right\} \stackrel{(1)}{\quad} \stackrel{(2)}{\quad} \Rightarrow \text{class}^r(h', \iota) = C$$

Proof.

$$\begin{aligned}
 \text{class}^r(h', \iota) &= (h[\iota_{tr} \mapsto \langle h(\iota_{tr}) \downarrow_1, \text{ow}, h(\iota_{tr}) \downarrow_3 \rangle]) \downarrow_1 \stackrel{(2)}{=} \\
 &= \left\{ \begin{array}{ll} \iota = \iota_{tr} \Rightarrow \langle h(\iota_{tr}) \downarrow_1, \text{ow}, h(\iota_{tr}) \downarrow_3 \rangle \downarrow_1 = h(\iota_{tr}) \downarrow_1 = \\ \iota \neq \iota_{tr} \Rightarrow h(\iota_{tr}) \downarrow_1 = \end{array} \right. \\
 &= \text{class}^r(h, \iota) \stackrel{(1)}{=} C
 \end{aligned}$$

□

Lemma 7.6.

$$\left. \begin{array}{l} h; \iota_0 \vdash_{wf} \iota : T \\ T \leq T' \end{array} \right\} \stackrel{(1)}{\quad} \stackrel{(2)}{\quad} \Rightarrow h; \iota_0 \vdash_{wf} \iota : T'$$

Proof.

$$\begin{aligned}
 T &= C \ m \\
 T' &= C' \ m' \\
 (2) &\Rightarrow \left\{ \begin{array}{ll} m \leq m' & (3) \\ C \sqsubseteq C' & (4) \end{array} \right. \\
 (3) &\Rightarrow \left\{ \begin{array}{ll} m = m' \stackrel{(1)}{\Rightarrow} h; \iota_0 \vdash_{wf} \iota : m' \\ m' = \text{any} \Rightarrow h; \iota_0 \vdash_{wf} \iota : \text{any} \end{array} \right. \\
 &\Rightarrow h; \iota_0 \vdash_{wf} \iota : m' \quad (5) \\
 (1) \wedge (4) &\Rightarrow C' \sqsubseteq \text{class}^r(h, \iota) \quad (6) \\
 (5) \wedge (6) &\Rightarrow h; \iota_0 \vdash_{wf} \iota : T'
 \end{aligned}$$

□

Lemma 7.7.

$$\left. \begin{array}{l} h; \iota_0 \vdash_{wf} \iota : T \\ h' = transferCl(h, \iota_{tr}, \text{ow}) \\ m \in \{\text{uniq}, \text{rep}\langle Cl \rangle\} \Rightarrow \text{owner}(h, \iota) \neq \text{owner}(h, \iota_{tr}) \\ \text{owner}(h, \iota_{tr}) \downarrow_2 = Cl \end{array} \right\} \stackrel{(1)}{\quad} \stackrel{(2)}{\quad} \stackrel{(3)}{\quad} \stackrel{(4)}{\quad} \Rightarrow h'; \iota_0 \vdash_{wf} \iota : T$$

Proof.

$$T = C \ m$$

$$(1) \Rightarrow C \sqsubseteq \text{class}^r(h, \iota) \xrightarrow{\text{Lem. 7.5}} C \sqsubseteq \text{class}^r(h', \iota) \quad (5)$$

$$m = \text{rep}\langle Cl' \rangle \xrightarrow{(1)} \text{owner}(h, \iota) = \langle Cl', \iota_0 \rangle \xrightarrow{(3) \wedge (4)} \text{owner}(h', \iota) = \langle Cl', \iota_0 \rangle \Rightarrow \\ \Rightarrow h'; \iota_0 \vdash_{wf} \iota : \text{rep}\langle Cl' \rangle \quad (6)$$

$$m = \text{peer} \xrightarrow{(1)} \text{owner}(h, \iota_0) = \text{owner}(h, \iota) \Rightarrow \\ \left\{ \begin{array}{l} \iota, \iota' \in \text{PeerSet}(h, \iota_{tr}) \Rightarrow \text{owner}(h', \iota_0) = \text{ow} \wedge \text{owner}(h', \iota) = \text{ow} \\ \iota, \iota' \notin \text{PeerSet}(h, \iota_{tr}) \Rightarrow \text{owner}(h', \iota_0) = \text{owner}(h', \iota) \end{array} \right. \Rightarrow \\ \Rightarrow \text{owner}(h', \iota_0) = \text{owner}(h', \iota) \Rightarrow h'; \iota_0 \vdash_{wf} \iota : \text{peer} \quad (7)$$

$$m = \text{uniq} \xrightarrow{(1)} \text{owner}(h, \iota) = \langle Cl, \text{null}^r \rangle \wedge Cl \in \text{CltFree} \xrightarrow{(3)} \\ \Rightarrow \text{owner}(h', \iota) = \langle Cl, \text{null}^r \rangle \wedge Cl \in \text{CltFree} \Rightarrow h'; \iota_0 \vdash_{wf} \iota : \text{uniq} \quad (8)$$

$$m = \text{this} \xrightarrow{(1)} \iota = \iota_0 \Rightarrow h'; \iota_0 \vdash_{wf} \iota : \text{this} \quad (9)$$

$$h'; \iota_0 \vdash_{wf} \iota : \text{any} \quad (10)$$

$$(6) \wedge (7) \wedge (8) \wedge (9) \wedge (10) \Rightarrow h'; \iota_0 \vdash_{wf} \iota : m \quad (11)$$

$$(5) \wedge (11) \Rightarrow h'; \iota_0 \vdash_{wf} \iota : T$$

□

Lemma 7.8.

$$\left. \begin{array}{l} h; \iota_0 \vdash_{wf} \iota : T \\ \text{owner}(h, \iota_0) = \text{owner}(h', \iota_0) \\ \text{owner}(h, \iota) = \text{owner}(h', \iota) \\ \text{class}^r(h, \iota) = \text{class}^r(h', \iota) \end{array} \right\} \stackrel{(1)}{\quad} \stackrel{(2)}{\quad} \stackrel{(3)}{\quad} \stackrel{(4)}{\quad} \Rightarrow h'; \iota_0 \vdash_{wf} \iota : T$$

Proof.

$$T = C \ m$$

$$(1) \Rightarrow C \sqsubseteq \text{class}^r(h, \ i) \xrightarrow{(4)} C \sqsubseteq \text{class}^r(h', \ i) \ (5)$$

$$\begin{aligned} m = \text{rep } \langle Cl \rangle &\xrightarrow{(1)} \text{owner}(h, \ i) = \langle Cl, \ i_0 \rangle \xrightarrow{(3)} \text{owner}(h', \ i) = \langle Cl, \ i_0 \rangle \Rightarrow \\ &\Rightarrow h'; i_0 \vdash_{wf} i : \text{rep } \langle Cl \rangle \ (6) \end{aligned}$$

$$\begin{aligned} m = \text{peer} &\xrightarrow{(1)} \text{owner}(h, \ i_0) = \text{owner}(h, \ i) \xrightarrow{(2) \wedge (3)} \\ &\Rightarrow \text{owner}(h', \ i_0) = \text{owner}(h', \ i) \Rightarrow h'; i_0 \vdash_{wf} i : \text{peer} \ (7) \end{aligned}$$

$$\begin{aligned} m = \text{uniq} &\xrightarrow{(1)} \text{owner}(h, \ i) = \langle Cl, \ \text{null}^r \rangle \wedge Cl \in \text{CltFree} \xrightarrow{(3)} \\ &\Rightarrow \text{owner}(h', \ i) = \langle Cl, \ \text{null}^r \rangle \wedge Cl \in \text{CltFree} \Rightarrow h'; i_0 \vdash_{wf} i : \text{uniq} \ (8) \end{aligned}$$

$$m = \text{this} \xrightarrow{(1)} i = i_0 \Rightarrow h'; i_0 \vdash_{wf} i : \text{this} \ (9)$$

$$h'; i_0 \vdash_{wf} i : \text{any} \ (10)$$

$$(6) \wedge (7) \wedge (8) \wedge (9) \wedge (10) \Rightarrow h'; i_0 \vdash_{wf} i : m \ (11)$$

$$(5) \wedge (11) \Rightarrow h'; i_0 \vdash_{wf} i : T$$

□

Lemma 7.9.

$$\left. \begin{array}{l} h; i_0 \vdash_{wf} i : T \\ \langle i', \ h' \rangle = \text{new}(h, \ C, \ ow) \end{array} \right\} \xrightarrow{(1) \atop (2)} h'; i_0 \vdash_{wf} i : T$$

Proof.

$$(1) \xrightarrow{(2) \wedge \text{Lem. 7.8}} h'; i_0 \vdash_{wf} i : T$$

□

Lemma 7.10.

$$\left. \begin{array}{l} h; i_0 \vdash_{wf} i : T \\ h' = h[i'_0 . f := i'] \end{array} \right\} \xrightarrow{(1) \atop (2)} h'; i_0 \vdash_{wf} i : T$$

Proof.

$$(1) \xrightarrow{(2) \wedge \text{Lem. 7.8}} h'; i_0 \vdash_{wf} i : T$$

□

7.3 heap's class well-formedness lemmas

Lemma 7.11.

$$\left. \begin{array}{l} \vdash_{wf} h \\ \iota_0 \in \text{dom}(h) \\ h' = h[\iota_0 . f' := \iota'] \\ h \vdash_{wf} \iota' : \text{class}(\text{fType}(\text{class}^r(\iota_0), f')) \end{array} \right\} \begin{array}{l} (1) \\ (2) \\ (3) \\ (4) \end{array} \Rightarrow \vdash_{wf} h'$$

Proof.

$$\begin{aligned} & \forall \iota \in \text{dom}(h') : C = \text{class}^r(h', \iota) \forall f \in \text{fields}(C) : \\ & \left\{ \begin{array}{l} \iota = \iota_0 \wedge f' = f \xrightarrow{(4)} \\ \iota \neq \iota_0 \vee f' \neq f \Rightarrow h'(\iota.f) = h(\iota.f) \xrightarrow{(1)} \\ \Rightarrow h' \vdash_{wf} h'(\iota.f) : \text{class}(\text{fType}(C, f)) \xrightarrow{(2)} \vdash_{wf} h' \end{array} \right. \end{aligned}$$

□

Lemma 7.12.

$$\left. \begin{array}{l} \vdash_{wf} h \\ \langle \iota', h' \rangle = \text{new}(h, C, \text{ow}) \end{array} \right\} \begin{array}{l} (1) \\ (2) \end{array} \Rightarrow \vdash_{wf} h'$$

Proof.

$$\begin{aligned} & \forall \iota \in \text{dom}(h') : C = \text{class}^r(h', \iota) \forall f \in \text{fields}(C) : \\ & \left\{ \begin{array}{l} \iota = \iota' \xrightarrow{(2)} h'(\iota.f) = \text{null}^r \Rightarrow \\ \iota \neq \iota' \Rightarrow h'(\iota.f) = h(\iota.f) \xrightarrow{(1)} \\ \Rightarrow h' \vdash_{wf} h'(\iota.f) : \text{class}(\text{fType}(C, f)) \xrightarrow{(2)} \vdash_{wf} h' \end{array} \right. \end{aligned}$$

□

Lemma 7.13.

$$\left. \begin{array}{l} \vdash_{wf} h \\ h' = \text{transferCl}(h, \iota_{tr}, \text{ow}) \end{array} \right\} \begin{array}{l} (1) \\ (2) \end{array} \Rightarrow \vdash_{wf} h'$$

Proof.

$$\begin{aligned} & \forall \iota \in \text{dom}(h') : C = \text{class}^r(h', \iota) \forall f \in \text{fields}(C) : \\ & \left\{ \begin{array}{l} \text{class}^r(h, \iota) = \text{class}^r(h', \iota) \\ h(\iota.f) = h'(\iota.f) = \iota' \\ \text{class}^r(h, \iota') = \text{class}^r(h', \iota') \end{array} \right. \Rightarrow \\ & \left. \begin{array}{l} \\ \\ \Rightarrow \\ h \vdash_{wf} h(\iota.f) : \text{class}(\text{fType}(C, f)) \\ \Rightarrow h' \vdash_{wf} h'(\iota.f) : \text{class}(\text{fType}(C, f)) \xrightarrow{(2)} \vdash_{wf} h' \end{array} \right. \end{aligned}$$

□

7.4 heap's well-formedness lemmas

Lemma 7.14.

$$\left. \begin{array}{l} U; \iota_0 \vdash_{wf} h \\ U' \cap \text{FieldID} \subseteq U \cap \text{FieldID} \end{array} \right\} \stackrel{(1)}{\Rightarrow} U'; \iota_0 \vdash_{wf} h$$

Proof.

$$\begin{aligned} & \forall \iota \in \text{dom}(h) : C = \text{class}^r(h, \iota) \forall f \in \text{fields}(C) : \\ & \langle \iota, \{\iota_0\} \times U \stackrel{(2)}{\Rightarrow} \langle \iota, \{\iota_0\} \times U' \stackrel{(1)}{\Rightarrow} h; \iota \vdash_{wf} h(\iota.f) : \text{mod(fType}(C, f)) \\ & \Rightarrow \text{FrSt}_2; h \vdash_{wf} \end{aligned}$$

□

Lemma 7.15.

$$\left. \begin{array}{l} \vdash_{wf} h \\ U; \iota_0 \vdash_{wf} h \\ h; \iota_0 \vdash_{wf} \iota'_0 : T \\ T = C.m \\ m = \text{this} \Rightarrow f \notin U \\ \iota'_0 \neq \text{null}^r \end{array} \right\} \stackrel{(1)}{\Rightarrow} h; \iota_0 \vdash_{wf} h(\iota'_0.f) : T \triangleright_U \text{fType}(C, f)$$

Proof.

$$\begin{aligned} (1) \wedge (2) \wedge (4) & \Rightarrow h; \iota'_0 \vdash_{wf} h(\iota'_0.f) : \text{fType}(C, f) \stackrel{(3) \wedge \text{Lem. 7.2}}{\Rightarrow} \\ & \Rightarrow h; \iota_0 \vdash_{wf} h(\iota'_0.f) : T \triangleright_U \text{fType}(C, f) \end{aligned}$$

□

Lemma 7.16.

$$\left. \begin{array}{l} U; \iota_0 \vdash_{wf} h \\ h; \iota'_0 \vdash_{wf} \iota' : \text{mod(fType(class}^r(\iota'_0), f)) \\ h' = h[\iota'_0.f' := \iota'] \\ h; \iota_0 \vdash_{wf} \iota'_0 : m \\ U' = \text{if}(m = \text{this}) \text{ then } U \setminus \{f\} \text{ else } U \end{array} \right\} \stackrel{(1)}{\Rightarrow} U'; \iota_0 \vdash_{wf} h'$$

Proof.

$$\begin{aligned} \forall \iota \in \text{dom}(h') : C = \text{class}^r(h', \iota) \forall f \in \text{fields}(C) : \\ \langle \iota, f \rangle \notin \{\iota_0\} \times U' \xrightarrow{(3) \wedge (4)} \\ \left\{ \begin{array}{l} \iota \neq \iota'_0 \vee f \neq f' \Rightarrow h(\iota.f) = h'(\iota.f) \xrightarrow{(1)} h'; \iota \vdash_{wf} h'(\iota.f) : \text{mod(fType}(C, f)) \\ \iota = \iota'_0 \wedge f = f' \xrightarrow{(2)} \\ \Rightarrow U'; \iota_0 \vdash_{wf} h' \end{array} \right. \end{aligned}$$

□

Lemma 7.17.

$$\left. \begin{array}{l} U; \iota_0 \vdash_{wf} h \\ U \cap \text{FieldID} = \emptyset \end{array} \right\} \xrightarrow{(1) \atop (2)} U'; \iota'_0 \vdash_{wf} h$$

Proof.

$$\begin{aligned} \forall \iota \in \text{dom}(h) : C = \text{class}^r(h, \iota) \forall f \in \text{fields}(C) : \\ \langle \iota, f \rangle \notin \{\iota'_0\} \times U' \Rightarrow \langle \iota, f \rangle \notin \{\iota_0\} \times U \xrightarrow{(1)} \\ h; \iota \vdash_{wf} h(\iota.f) : \text{mod(fType}(C, f)) \\ \Rightarrow U'; \iota'_0 \vdash_{wf} h \end{aligned}$$

□

Lemma 7.18.

$$\left. \begin{array}{l} U; \iota_0 \vdash_{wf} h \\ \langle \iota', h' \rangle = \text{new}(h, C, \text{ow}) \end{array} \right\} \xrightarrow{(1) \atop (2)} U; \iota_0 \vdash_{wf} h'$$

Proof.

$$\begin{aligned} \forall \iota \in \text{dom}(h') : C = \text{class}^r(h', \iota) \forall f \in \text{fields}(C) : \\ \langle \iota, f \rangle \notin \{\iota_0\} \times U \xrightarrow{(2)} \\ \left\{ \begin{array}{l} \iota \neq \iota' \Rightarrow h(\iota.f) = h'(\iota.f) \xrightarrow{(1)} h'; \iota \vdash_{wf} h'(\iota.f) : \text{mod(fType}(C, f)) \\ \iota = \iota' \Rightarrow h'(\iota.f) = \text{null}^r \\ \Rightarrow U; \iota_0 \vdash_{wf} h' \end{array} \right. \end{aligned}$$

□

Lemma 7.19.

$$\left. \begin{array}{l} \Gamma(y) = \mathbf{rep}\langle Cl \rangle \\ h; \iota_0 \vdash_{wf} \iota_{tr} : \mathbf{rep}\langle Cl \rangle \\ h' = \mathbf{transferCl}(h, \iota_{tr}, \mathbf{ow}) \\ U; \iota_0 \vdash_{wf} h \\ U' = \mathbf{consumeAliases}(U, y) \end{array} \right\} \begin{array}{l} (1) \\ (2) \\ (3) \\ (4) \\ (5) \end{array} \Rightarrow U'; \iota_0 \vdash_{wf} h'$$

Proof.

$$\begin{aligned} & \forall \iota \in \text{dom}(h') : C = \mathbf{class}^r(h', \iota) \forall f \in \mathbf{fields}(C) : \\ & \langle \iota, f \rangle \notin \{\iota_0\} \times U' \stackrel{(1) \wedge (2) \wedge (3) \wedge (5)}{\Rightarrow} \\ & \mathbf{mod}(\mathbf{fType}(C, f)) \neq \mathbf{rep}\langle Cl \rangle \vee \iota \neq \iota_0 \stackrel{(4) \wedge \text{Lem. 7.7}}{\Rightarrow} \\ & h'; \iota \vdash_{wf} h'(\iota.f) : \mathbf{mod}(\mathbf{fType}(C, f)) \\ & \Rightarrow U'; \iota_0 \vdash_{wf} h' \end{aligned}$$

□

Lemma 7.20.

$$\left. \begin{array}{l} h; \iota_0 \vdash_{wf} \iota_{tr} : \mathbf{uniq} \\ h' = \mathbf{transferCl}(h, \iota_{tr}, \mathbf{ow}) \\ U; \iota_0 \vdash_{wf} h \end{array} \right\} \begin{array}{l} (1) \\ (2) \\ (3) \end{array} \Rightarrow U; \iota_0 \vdash_{wf} h'$$

Proof.

$$\begin{aligned} & \forall \iota \in \text{dom}(h') : C = \mathbf{class}^r(h', \iota) \forall f \in \mathbf{fields}(C) : \\ & \langle \iota, f \rangle \notin \{\iota_0\} \times U \stackrel{(1) \wedge (3)}{\Rightarrow} \\ & \mathbf{mod}(\mathbf{fType}(C, f)) \neq \mathbf{rep}\langle Cl \rangle \vee \mathbf{owner}(h, \iota) \neq \mathbf{owner}(h, \iota_{tr}) \stackrel{(2) \wedge \text{Lem. 7.7}}{\Rightarrow} \\ & h'; \iota \vdash_{wf} h'(\iota.f) : \mathbf{mod}(\mathbf{fType}(C, f)) \\ & \Rightarrow U; \iota_0 \vdash_{wf} h' \end{aligned}$$

□

7.5 Frames stack's well-formedness lemmas

Lemma 7.21.

$$\left. \begin{array}{l} h \vdash_{wf} \Gamma^r; \Gamma; U \\ U' \cap \mathbf{VarID} \subseteq U \cap \mathbf{VarID} \end{array} \right\} \begin{array}{l} (1) \\ (2) \end{array} \Rightarrow h \vdash_{wf} \Gamma^r; \Gamma; U'$$

Proof.

$$(1) \Rightarrow \left\{ \begin{array}{ll} \text{dom}(\Gamma^r) = \text{dom}(\Gamma) & (3) \\ \text{dom}(\Gamma) \supseteq U & (4) \\ \text{mod}(\Gamma(\text{this})) = \text{this} & (5) \\ \forall v \in \text{dom}(\Gamma^r) \setminus U : h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^r(v) : \Gamma(v) & (6) \end{array} \right. \\ (2) \wedge (6) \Rightarrow \forall v \in \text{dom}(\Gamma^r) \setminus U' : h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^r(v) : \Gamma(v) \quad (7) \\ (3) \wedge (4) \wedge (5) \wedge (7) \Rightarrow h \vdash_{wf} \Gamma^r; \Gamma; U'$$

□

Lemma 7.22.

$$\left. \begin{array}{ll} h \vdash_{wf} \Gamma^r; \Gamma; U & (1) \\ U' = U \cup \{x\} & (2) \\ \Gamma^{r'} = \Gamma^r[x \mapsto \iota'] & (3) \\ h; \Gamma^r(\text{this}) \vdash_{wf} \iota' : \Gamma(x) & (4) \\ x \in \text{dom}(\Gamma) & (5) \end{array} \right\} \Rightarrow h \vdash_{wf} \Gamma^{r'}; \Gamma; U'$$

Proof.

$$\left. \begin{array}{ll} \text{dom}(\Gamma^r) = \text{dom}(\Gamma) & (6) \\ \text{dom}(\Gamma) \supseteq U & (7) \\ \text{mod}(\Gamma(\text{this})) = \text{this} & (8) \\ \forall v \in \text{dom}(\Gamma^r) \setminus U : h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^r(v) : \Gamma(v) & (9) \\ v \in \text{dom}(\Gamma^{r'}) \setminus U' \left\{ \begin{array}{ll} v = x \xrightarrow{(4)} h; \Gamma^{r'}(\text{this}) \vdash_{wf} \Gamma^{r'}(v) : \Gamma(v) & (10) \\ v \neq x \xrightarrow{(9)} h; \Gamma^{r'}(\text{this}) \vdash_{wf} \Gamma^{r'}(v) : \Gamma(v) & (10) \end{array} \right. \\ (5) \wedge (6) \Rightarrow \text{dom}(\Gamma^{r'}) = \text{dom}(\Gamma) & (11) \\ (11) \wedge (7) \wedge (8) \wedge (10) \Rightarrow h \vdash_{wf} \Gamma^{r'}; \Gamma; U' \end{array} \right.$$

□

Lemma 7.23.

$$\left. \begin{array}{ll} h \vdash_{wf} \Gamma^r; \Gamma; U & (1) \\ \iota_0 \in \text{dom}(h) & (2) \\ h' = h[\iota_0 . f := \iota] & (3) \end{array} \right\} \Rightarrow h' \vdash_{wf} \Gamma^r; \Gamma; U$$

Proof.

$$\begin{aligned}
 (1) &\Rightarrow \left\{ \begin{array}{ll} \text{dom}(\Gamma^r) = \text{dom}(\Gamma) & (4) \\ \text{dom}(\Gamma) \supseteq U & (5) \\ \text{mod}(\Gamma(\text{this})) = \text{this} & (6) \\ \forall v \in \text{dom}(\Gamma^r) \setminus U : h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^r(v) : \Gamma(v) & (7) \end{array} \right. \\
 (7) &\Rightarrow \forall v \in \text{dom}(\Gamma^r) \setminus U : h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^r(v) : \Gamma(v) \\
 (3) \wedge \text{Lem. } 7.10 &\Rightarrow h'; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^r(v) : \Gamma(v) \quad (8) \\
 (4) \wedge (5) \wedge (6) \wedge (8) &\Rightarrow h' \vdash_{wf} \Gamma^r; \Gamma; U
 \end{aligned}$$

□

Lemma 7.24.

$$\left. \begin{array}{l} h \vdash_{wf} \Gamma^r; \Gamma; U \\ \Gamma^{r'} = \{\langle p, \Gamma^r(z) \rangle, \overline{\langle v, \text{null}^r \rangle}, \langle \text{this}, \Gamma^r(y) \rangle, \langle \text{res}, \text{null}^r \rangle\} \\ \Gamma' = \{\langle p, T_p \rangle, \overline{\langle v, T \rangle}, \langle \text{this}, \text{this class}^r(\Gamma^r(y)) \rangle, \langle \text{res}, T_r \rangle\} \\ \Gamma(z) \leq \Gamma(y) \triangleright_U T_p \end{array} \right\} \begin{array}{l} (1) \\ (2) \\ (3) \\ (4) \end{array} \Rightarrow h \vdash_{wf} \Gamma^{r'}; \Gamma'; \emptyset$$

Proof.

$$\begin{aligned}
 \text{dom}(\Gamma^{r'}) &= \{p, \bar{v}, \text{this}, \text{res}\} = \text{dom}(\Gamma') \Rightarrow \text{dom}(\Gamma^{r'}) = \text{dom}(\Gamma') \quad (5) \\
 \text{dom}(\Gamma') &\supseteq \emptyset \quad (6) \\
 \text{mod}(\Gamma'(\text{this})) &= \text{this} \quad (7) \\
 (1) &\Rightarrow h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^{r'}(\text{this}) : \Gamma(y) \wedge h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^{r'}(p) : \Gamma(z) \wedge (4) \\
 \text{Lem. } 7.3 &\Rightarrow h; \Gamma^{r'}(\text{this}) \vdash_{wf} \Gamma^{r'}(p) : T_p \quad (8) \\
 \Gamma^{r'}(\text{this}) &= \Gamma^r(\text{this}) \Rightarrow h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^{r'}(\text{this}) : \Gamma'(\text{this}) \quad (9) \\
 \Gamma^{r'}(\bar{v}) &= \text{null}^r \Rightarrow h; \Gamma^{r'}(\bar{v}) \vdash_{wf} \Gamma^{r'}(\bar{v}) : \Gamma'(\bar{v}) \quad (10) \\
 \Gamma^{r'}(\text{res}) &= \text{null}^r \Rightarrow h; \Gamma^{r'}(\text{res}) \vdash_{wf} \Gamma^{r'}(\text{res}) : \Gamma'(\text{res}) \quad (11) \\
 (8) \wedge (9) \wedge (10) \wedge (11) &\Rightarrow \forall v \in \text{dom}(\Gamma^{r'}) \setminus \emptyset : h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^{r'}(v) : \Gamma'(v) \quad (12) \\
 (5) \wedge (6) \wedge (7) \wedge (12) &\Rightarrow h \vdash_{wf} \Gamma^{r'}; \Gamma'; \emptyset
 \end{aligned}$$

□

Lemma 7.25.

$$\left. \begin{array}{l} FrSt = Fr \circ FrSt' \\ h \vdash_{wf} FrSt \end{array} \right\} \begin{array}{l} (1) \end{array} \Rightarrow h \vdash_{wf} FrSt' h$$

Proof.

$$\text{Fr} \circ \text{FrSt}' = \langle \Gamma^r_0, \Gamma^r_0, \Gamma^r_0 \rangle \circ \overline{\langle \Gamma^r, \Gamma, U \rangle}$$

$$(1) \Rightarrow h_0 \vdash_{wf} \Gamma^r_0; \Gamma^r_0; \Gamma^r_0 \wedge \overline{h \vdash_{wf} \Gamma^r; \Gamma^r; \Gamma^r} \Rightarrow \overline{h \vdash_{wf} \Gamma^r; \Gamma^r; \Gamma^r} \Rightarrow h \vdash_{wf} \text{FrSt}' h$$

□

Lemma 7.26.

$$h \vdash_{wf} \Gamma^r; \Gamma; U$$

$$\langle \iota, h' \rangle = \text{new}(h, C, \text{ow}) \quad \begin{cases} (1) \\ (2) \end{cases} \quad \left\{ \begin{array}{l} \Rightarrow h' \vdash_{wf} \Gamma^r; \Gamma; U \end{array} \right.$$

Proof.

$$(1) \Rightarrow \begin{cases} \text{dom}(\Gamma^r) = \text{dom}(\Gamma) & (4) \\ \text{dom}(\Gamma) \supseteq U & (5) \\ \text{mod}(\Gamma(\text{this})) = \text{this} & (6) \\ \forall v \in \text{dom}(\Gamma^r) \setminus U : h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^r(v) : \Gamma(v) & (7) \end{cases}$$

$$(7) \Rightarrow \forall v \in \text{dom}(\Gamma^r) \setminus U : h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^r(v) : \Gamma(v)$$

$$\stackrel{(4) \wedge \text{Lem. 7.9}}{\Rightarrow} h'; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^r(v) : \Gamma(v) \quad (8)$$

$$(4) \wedge (5) \wedge (6) \wedge (8) \Rightarrow h' \vdash_{wf} \Gamma^r; \Gamma; U$$

□

Lemma 7.27.

$$\left. \begin{array}{l} \text{Fr} = \langle \Gamma^r, \Gamma, U \rangle \\ \text{FrSt} = \text{Fr} \circ \text{FrSt}_0 \\ \text{Fr}' = \langle \Gamma'^r, \Gamma, U' \rangle \\ \text{FrSt}' = \text{Fr}' \circ \text{FrSt}_0 \\ \Gamma(y) = \text{rep} \langle Cl \rangle & (1) \\ \Gamma(x) = \text{uniq} & (2) \\ h \vdash_{wf} \text{FrSt} & (3) \\ h \vdash_{Inv} \text{FrSt} & (4) \\ \text{ow} = \langle \text{newClt}(h), \text{null}^r \rangle & (5) \\ h' = \text{transferCl}(h, \Gamma^r(y), \text{ow}) & (6) \\ U' = \text{consumeAliases}(U, y) \setminus \{x\} & (7) \\ \Gamma'^r = \Gamma^r[x \mapsto \Gamma^r(y)] & (8) \end{array} \right\} \Rightarrow h' \vdash_{wf} \text{FrSt}'$$

Proof.

$$\begin{aligned}
 & (3) \wedge (4) \wedge (6) \xrightarrow{(4) \wedge \text{Lem. 7.7}} h' \vdash_{wf} \text{FrSt}_0 \quad (9) \\
 & (3) \Rightarrow \left\{ \begin{array}{l} \text{dom}(\Gamma^r) = \text{dom}(\Gamma) \\ \text{dom}(\Gamma) \supseteq U \\ \text{mod}(\Gamma(\text{this})) = \text{this} \\ \forall v \in \text{dom}(\Gamma^r) \setminus U : h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^r(v) : \Gamma(v) \end{array} \right. \quad (10) \\
 & \forall v' \in \text{dom}(\Gamma^{r'}) \setminus U' \xrightarrow{(7) \wedge (8)} \Gamma^{r'}(v) \neq \text{rep} \langle Cl \rangle \wedge v \neq y \xrightarrow{\text{Lem. 7.7} \wedge (3)} \\
 & h'; \Gamma^{r'}(\text{this}) \vdash_{wf} \Gamma^{r'}(v) : \Gamma'(v) \quad (14) \\
 & (10) \wedge (11) \wedge (12) \wedge (14) \Rightarrow h' \vdash_{wf} \Gamma^r; \Gamma; U \quad (15) \\
 & (9) \wedge (15) \Rightarrow h' \vdash_{wf} \text{FrSt}'
 \end{aligned}$$

□

Lemma 7.28.

$$\left. \begin{array}{ll} Fr = \langle \Gamma^r, \Gamma, U \rangle \\ FrSt = Fr \circ FrSt_0 \\ Fr' = \langle \Gamma^{r'}, \Gamma, U' \rangle \\ FrSt' = Fr' \circ FrSt_0 \\ \Gamma(y) = \text{uniq} & (1) \\ \Gamma(x) = \text{rep} \langle Cl \rangle & (2) \\ h \vdash_{wf} FrSt & (3) \\ FrSt \vdash_{Uniq} h & (4) \\ ow = \text{mod20w}(h, m, \Gamma^r(\text{this})) & (5) \\ m \in \{\text{rep} \langle Cl \rangle, \text{peer}\} & (6) \\ h' = \text{transferCl}(h, \Gamma^r(y), ow) & (7) \\ U' = U \cup \{y\} \setminus \{x\} & (8) \\ \Gamma^{r'} = \Gamma^r[x \mapsto \Gamma^r(y)] & (9) \end{array} \right\} \Rightarrow h' \vdash_{wf} FrSt'$$

Proof.

$$\begin{aligned}
 & (3) \wedge (4) \wedge (6) \xrightarrow{(4) \wedge \text{Lem. 7.7}} h' \vdash_{wf} \text{FrSt}_0 \quad (9) \\
 & (3) \Rightarrow \left\{ \begin{array}{l} \text{dom}(\Gamma^r) = \text{dom}(\Gamma) \\ \text{dom}(\Gamma) \supseteq U \\ \text{mod}(\Gamma(\text{this})) = \text{this} \\ \forall v \in \text{dom}(\Gamma^r) \setminus U : h; \Gamma^r(\text{this}) \vdash_{wf} \Gamma^r(v) : \Gamma(v) \end{array} \right. \quad (10) \\
 & \forall v' \in \text{dom}(\Gamma^{r'}) \setminus U' \xrightarrow{(7) \wedge (8)} v \neq y \xrightarrow{\text{Lem. 7.7} \wedge (3)} \\
 & h'; \Gamma^{r'}(\text{this}) \vdash_{wf} \Gamma^{r'}(v) : \Gamma'(v) \quad (14) \\
 & (10) \wedge (11) \wedge (12) \wedge (14) \Rightarrow h' \vdash_{wf} \Gamma^r; \Gamma; U \quad (15) \\
 & (9) \wedge (15) \Rightarrow h' \vdash_{wf} \text{FrSt}'
 \end{aligned}$$

□

7.6 Uniqueness's lemmas

Lemma 7.29.

$$\left. \begin{array}{l} FrSt = \langle \Gamma^r, \Gamma, U \rangle \circ FrSt_0 \\ h \vdash_{Uniq} FrSt \\ U' \cap VarID \subseteq U \cap VarID \\ FrSt' = \langle \Gamma^r, \Gamma, U' \rangle \circ FrSt_0 \end{array} \right\} \begin{array}{l} (1) \\ (2) \\ (3) \end{array} \Rightarrow h \vdash_{Uniq} FrSt'$$

Proof.

$$(1) \Rightarrow \forall Cl \in \text{CltFree} \text{ RefNum}(h, FrSt, \langle Cl, \text{null}^r \rangle, \text{uniq}) \leq 1 \xrightarrow{(2)} \\ \Rightarrow \text{RefNum}(h, FrSt', \langle Cl, \text{null}^r \rangle, \text{uniq}) \leq 1 \Rightarrow h \vdash_{Uniq} FrSt'$$

□

Lemma 7.30.

$$\left. \begin{array}{l} FrSt = \langle \Gamma^r, \Gamma, U \rangle \circ FrSt_0 \\ h \vdash_{Uniq} FrSt \\ U' = U \cup \{x\} \\ \Gamma^{r'} = \Gamma^r[x \mapsto \iota'] \\ \iota' \neq \text{null}^r \wedge \Gamma(x) = \text{uniq} \Rightarrow \\ \text{RefNum}(h, FrSt, \text{owner}(h, \iota'), \text{uniq}) = 0 \\ FrSt' = \langle \Gamma^{r'}, \Gamma, U' \rangle \circ FrSt_0 \end{array} \right\} \begin{array}{l} (1) \\ (2) \\ (3) \\ (4) \end{array} \Rightarrow h \vdash_{Uniq} FrSt'$$

Proof.

$$(1) \Rightarrow \forall Cl \in \text{CltFree} \left\{ \begin{array}{l} Cl = \text{owner}(h, \iota') \downarrow_1 \xrightarrow{(3) \wedge (4)} \\ Cl \neq \text{owner}(h, \iota') \downarrow_1 \xrightarrow{(1)} \end{array} \right. \\ \text{RefNum}(h, FrSt', \langle Cl, \text{null}^r \rangle, \text{uniq}) \leq 1 \Rightarrow h \vdash_{Uniq} FrSt'$$

□

Lemma 7.31.

$$\left. \begin{array}{l} h \vdash_{Uniq} FrSt \\ \iota_0 \in \text{dom}(h) \\ h' = h[\iota_0 . f := \iota] \end{array} \right\} \begin{array}{l} (1) \\ (2) \\ (3) \end{array} \Rightarrow h' \vdash_{Uniq} FrSt$$

Proof.

$$(1) \Rightarrow \forall Cl \in \text{CltFree} \text{ RefNum}(h, \text{FrSt}, \langle Cl, \text{null}^r \rangle, \text{uniq}) \leq 1 \stackrel{(2) \wedge (3)}{\Rightarrow} \\ \Rightarrow \text{RefNum}(h', \text{FrSt}, \langle Cl, \text{null}^r \rangle, \text{uniq}) \leq 1 \Rightarrow h' \vdash_{\text{Uniq}} \text{FrSt}$$

1

Lemma 7.32.

$$\left. \begin{array}{l} FrSt = \langle \Gamma^r, \Gamma, U \rangle \circ FrSt_0 \\ h \vdash_{Uniq} FrSt \\ \Gamma^{r'} = \{\langle p, \Gamma^r(z) \rangle, \overline{\langle v, null^r \rangle}, \langle this, \Gamma^r(y) \rangle, \langle res, null^r \rangle\} \\ \Gamma' = \{\langle p, T_p \rangle, \overline{\langle v, T \rangle}, \langle this, this\ class^r(\Gamma^r(y)) \rangle, \langle res, T_r \rangle\} \\ U' = consumeUniq(U, T_p, z) \\ Fr_1 = \langle \Gamma^{r'}, \Gamma', \emptyset \rangle \\ Fr_2 = \langle \Gamma^r, \Gamma, U' \rangle \\ \Rightarrow h \vdash_{Uniq} Fr_1 \circ Fr_2 \circ FrSt_0 \end{array} \right\} \Rightarrow$$

Proof.

$$\begin{aligned}
(1) \Rightarrow & \forall Cl \in \text{CltFree} \setminus \{\text{owner}(h, \Gamma^{r'}(p)) \downarrow_1\} \Rightarrow \\
& \Rightarrow \text{RefNum}(h, Fr_1 \circ Fr_2 \circ FrSt_0, \langle Cl, \text{null}^r \rangle, \text{uniq}) \leq 1 \quad (5) \\
\Gamma'(p) \neq \text{uniq} \Rightarrow & \text{owner}(h, \Gamma^{r'}(p)) \downarrow_1 \notin Cl \quad (6) \\
\Gamma'(p) = \text{uniq} \Rightarrow & \text{RefNum}(h, Fr_1 \circ Fr_2 \circ FrSt_0, \langle Cl, \text{null}^r \rangle, \text{uniq}) = \\
= & \text{RefNum}(h, FrSt, \langle Cl, \text{null}^r \rangle, \text{uniq}) \leq 1 \quad (7) \\
(5) \wedge (6) \wedge (7) \Rightarrow & h \vdash_{\text{Uniq}} Fr_1 \circ Fr_2 \circ FrSt_0
\end{aligned}$$

1

Lemma 7.33.

$$\left. \begin{array}{l} FrSt = Fr \circ FrSt' \\ h \vdash_{Uniq} FrSt \end{array} \right\} (1) \Rightarrow h \vdash_{Uniq} FrSt'$$

Proof.

$$\begin{aligned} \forall Cl \in \text{CltFree}(1) \Rightarrow \text{RefNum}(h, \text{FrSt}', \langle Cl, \text{null}^r \rangle, \text{uniq}) &\leq \\ \leq \text{RefNum}(h, \text{FrSt}, \langle Cl, \text{null}^r \rangle, \text{uniq}) &\leq 1 \Rightarrow h' \vdash_{\text{Uniq}} \text{FrSt} \end{aligned}$$

1

Lemma 7.34.

$$\left. \begin{array}{l} h \vdash_{Uniq} FrSt \\ \langle \iota, h' \rangle = new(h, C, ow) \end{array} \right\} \stackrel{(1)}{(2)} \Rightarrow h' \vdash_{Uniq} FrSt$$

Proof.

$$\begin{aligned} \forall Cl \in \text{CltFree}(1) \Rightarrow \text{RefNum}(h, \text{FrSt}', \langle Cl, \text{null}^r \rangle, \text{uniq}) = \\ = \text{RefNum}(h, \text{FrSt}, \langle Cl, \text{null}^r \rangle, \text{uniq}) \leq 1 \Rightarrow h' \vdash_{\text{Uniq}} \text{FrSt} \end{aligned}$$

□

Lemma 7.35.

$$\left. \begin{array}{l} Fr = \langle \Gamma^r, \Gamma, U \rangle \\ FrSt = Fr \circ FrSt_0 \\ Fr' = \langle \Gamma^{r'}, \Gamma, U' \rangle \\ FrSt' = Fr' \circ FrSt_0 \\ \Gamma(y) = rep \langle Cl \rangle \\ h \vdash_{wf} FrSt \\ FrSt \vdash_{Uniq} h \\ ow = \langle null^r, newClt(h) \rangle \\ h' = transferCl(h, \Gamma^r(y), ow) \\ U' = consumeAliases(U, y) \setminus \{x\} \\ \Gamma^{r'} = \Gamma^r[x \mapsto \Gamma^r(y)] \end{array} \right\} \Rightarrow h' \vdash_{Uniq} FrSt'$$

Proof.

$$\begin{aligned} & \forall Cl \in \text{CltFree} \setminus \{\text{ow } \downarrow_1\} \text{RefNum}(h', \text{FrSt}', \langle Cl, \text{null}^r \rangle, \text{uniq}) \stackrel{(2) \wedge (5) \wedge (6)}{\leq} \\ & \leq \text{RefNum}(h, \text{FrSt}, \langle Cl, \text{null}^r \rangle, \text{uniq}) \leq 1 \quad (8) \\ (4) \Rightarrow & \text{RefNum}(h, \text{FrSt}, \langle \text{ow } \downarrow_1, \text{null}^r \rangle, \text{uniq}) = 0 \Rightarrow \\ \Rightarrow & \text{RefNum}(h', \text{FrSt}', \langle \text{ow } \downarrow_1, \text{null}^r \rangle, \text{uniq}) = 1 \quad (9) \\ (8) \wedge (9) \Rightarrow & h' \vdash_{\text{Uniq}} \text{FrSt}' \end{aligned}$$

1

Lemma 7.36.

$$\left. \begin{array}{ll} Fr = \langle \Gamma^r, \Gamma, U \rangle \\ FrSt = Fr \circ FrSt_0 \\ Fr' = \langle \Gamma'^r, \Gamma, U' \rangle \\ FrSt' = Fr' \circ FrSt_0 \\ \Gamma(y) = uniq & (1) \\ h \vdash_{wf} FrSt & (2) \\ FrSt \vdash_{Uniq} h & (3) \\ ow \neq \langle null^r, _ \rangle & (4) \\ h' = transferCl(h, \Gamma^r(y), ow) & (5) \\ U' = U \cup \{y\} \setminus \{x\} & (6) \\ \Gamma'^r = \Gamma^r[x \mapsto \Gamma^r(y)] & (7) \end{array} \right\} \Rightarrow h' \vdash_{Uniq} FrSt'$$

Proof.

$$\begin{aligned} \forall Cl \in \text{CltFree} \quad (3) \Rightarrow \text{RefNum}(h', FrSt', \langle Cl, \text{null}^r \rangle, \text{uniq}) &\stackrel{(2) \wedge (5) \wedge (6)}{\Rightarrow} \\ \leq \text{RefNum}(h, FrSt, \langle Cl, \text{null}^r \rangle, \text{uniq}) \leq 1 \Rightarrow h' \vdash_{Uniq} FrSt' \end{aligned}$$

□

7.7 Global invariant's lemmas

Lemma 7.37.

$$\left. \begin{array}{ll} Fr = \langle \Gamma^r, \Gamma, U \rangle \\ Fr' = \langle \Gamma'^r, \Gamma, U' \rangle \\ FrSt = Fr \circ FrSt_0 \\ FrSt' = Fr' \circ FrSt_0 \\ h \vdash_{Inv} FrSt & (1) \\ \Gamma^r(\text{this}) = \Gamma'^r(\text{this}) & (2) \\ h \vdash_{wf} FrSt' & (3) \\ \forall \iota : ExtOwn(h, Fr, \iota) \Rightarrow ExtOwn(h, Fr', \iota) & (4) \end{array} \right\} \Rightarrow h \vdash_{Inv} FrSt'$$

Proof.

$$\begin{aligned} \iota_0 &= \Gamma^r(\text{this}) \\ (1) \Rightarrow & \\ \left\{ \begin{array}{l} h; FrSt' \vdash_{Inv} \iota_0 \\ \forall \iota : ExtOwn(h, Fr, \iota) \wedge \iota \neq \iota_0 \Rightarrow h; FrSt \vdash_{Inv} \iota \end{array} \right. & (5) \\ (6) \\ (4) \wedge (3) \wedge (6) \Rightarrow \forall \iota : ExtOwn(h', Fr', \iota) \wedge \iota \neq \iota_0 \Rightarrow h'; FrSt' \vdash_{Inv} \iota & (7) \\ (5) \wedge (7) \Rightarrow h \vdash_{Inv} FrSt' \end{aligned}$$

□

Lemma 7.38.

$$\left. \begin{array}{l} h \vdash_{Inv} FrSt \\ \iota_0 \in \text{dom}(h) \\ h' = h[\iota_0.f := \iota] \end{array} \right\} \stackrel{(1)}{\quad} \stackrel{(2)}{\quad} \stackrel{(3)}{\quad} \Rightarrow h' \vdash_{Inv} FrSt$$

Proof.

$$(2) \wedge (3) \Rightarrow \forall \iota \in \text{dom}(h') : \text{owner}(h, \iota) = \text{owner}(h', \iota) \xrightarrow{(2)} h' \vdash_{Inv} FrSt$$

□

Lemma 7.39.

$$\left. \begin{array}{l} Fr = \langle \Gamma^r, \Gamma, U \rangle \\ FrSt = Fr \circ FrSt_0 \\ \Gamma^{r'} = \{ \langle p, \Gamma^r(z) \rangle, \overline{\langle v, \text{null}^r \rangle}, \langle \text{this}, \Gamma^r(y) \rangle, \langle \text{res}, \text{null}^r \rangle \} \\ \Gamma' = \{ \langle p, T_p \rangle, \overline{\langle v, T \rangle}, \langle \text{this}, \text{this class}^r(\Gamma^r(y)) \rangle, \langle \text{res}, T_r \rangle \} \\ U' = \text{if } (\Gamma(y) \notin \{\text{peer, this}\}) \text{ then } U \text{ else } \text{consumeLocals}(U) \\ Fr_1 = \langle \Gamma^{r'}, \Gamma', \emptyset \rangle \\ Fr_2 = \langle \Gamma^r, \Gamma, U' \rangle \\ FrSt' = Fr_1 \circ Fr_2 \circ FrSt_0 \\ h \vdash_{Inv} FrSt \\ h \vdash_{wf} FrSt \\ h \vdash_{wf} FrSt' \\ \text{IsWritable}(\text{class}^r(\Gamma^r(y))) \\ \Rightarrow h \vdash_{Inv} FrSt' \end{array} \right\} \Rightarrow \left. \begin{array}{l} (1) \\ (2) \\ (3) \\ (4) \end{array} \right\}$$

Proof.

$$\begin{aligned} \iota_0 &= \Gamma^r(\text{this}) \\ (1) &\Rightarrow \\ \left\{ \begin{array}{l} h; FrSt_0 \vdash_{Inv} \iota_0 \\ \forall \iota : \text{ExtOwn}(h, Fr, \iota) \wedge \iota \neq \iota_0 \Rightarrow h; FrSt \vdash_{Inv} \iota \end{array} \right. & (5) \\ \iota'_0 &= \Gamma^{r'}(\text{this}) \\ (2) \wedge (3) \wedge (4) &\Rightarrow (\forall \iota : \text{ExtOwn}(h, Fr, \iota) \Rightarrow \text{ExtOwn}(h, Fr', \iota)) \Rightarrow \\ &\text{ExtOwn}(h', Fr', \iota) \wedge \iota \neq \iota'_0 \Rightarrow h'; Fr \circ FrSt_0 \vdash_{Inv} \iota & (7) \\ \Gamma(y) &\in \{\text{peer, this}\} \Rightarrow U' = \text{consumeLocals}(U) \Rightarrow \\ &\stackrel{(1)}{\Rightarrow} h; FrSt_1 \vdash_{Inv} \iota'_0 & (8) \\ \Gamma(y) = \text{rep } \langle Cl \rangle &\xrightarrow{(1)} h; FrSt_1 \vdash_{Inv} \iota'_0 & (9) \\ (7) \wedge (8) \wedge (9) &\Rightarrow h \vdash_{Inv} FrSt' \end{aligned}$$

1

Lemma 7.40.

$$\left. \begin{array}{l} Fr_1 = \langle \Gamma^r_1, \Gamma_1, U_1 \rangle \\ Fr_2 = \langle \Gamma^r_2, \Gamma_2, U_2 \rangle \\ Fr_3 = \langle \Gamma^r_3, \Gamma_3, U_3 \rangle \\ FrSt_1 = Fr_1 \circ FrSt_0 \\ FrSt_3 = Fr_3 \circ FrSt_1 \\ h; \Gamma^r_1(this) \vdash_{wf} \Gamma^r_2(this) : rep \langle Cl \rangle \quad (1) \\ h \vdash_{Inv} FrSt_1 \quad (2) \\ h' \vdash_{Inv} FrSt_3 \quad (3) \\ \Gamma^r_2(this) = \Gamma^r_3(this) \quad (4) \\ OwAsMod(h, h', Fr_2, Fr_3) \quad (5) \end{array} \right\} \Rightarrow h' \vdash_{Inv} FrSt_1$$

Proof.

$$\forall \iota : \text{ExtOwn}(h', \text{Fr}_1, \Gamma^r_1(\text{this})) \stackrel{(1)}{\Rightarrow} \begin{cases} \text{ExtOwn}(h', \text{Fr}_2, \Gamma^r_3(\text{this})) \stackrel{(2) \wedge (5)}{\Rightarrow} \\ \text{ExtOwn}(h', \text{Fr}_3, \Gamma^r_3(\text{this})) \stackrel{(3)}{\Rightarrow} \end{cases}$$

$$h'; \text{FrSt}_1 \vdash_{\text{Inv}} \iota \Rightarrow h' \vdash_{\text{Inv}} \text{FrSt}_1$$

1

Lemma 7.41.

$$\left. \begin{array}{l} Fr_1 = \langle \Gamma^r_1, \Gamma_1, U_1 \rangle \\ Fr_2 = \langle \Gamma^r_2, \Gamma_2, U_2 \rangle \\ FrSt_1 = Fr_1 \circ FrSt_0 \\ FrSt_2 = Fr_2 \circ FrSt_1 \\ h; \Gamma^r_1(this) \vdash_{wf} \Gamma^r_2(this) : peer \quad (1) \\ h \vdash_{Inv} FrSt_2 \quad (2) \end{array} \right\} \Rightarrow h \vdash_{Inv} FrSt_1$$

Proof.

$$\begin{aligned} \forall \ell : \text{ExtOwn}(h, Fr_1, \Gamma^r_2(\text{this})) &\stackrel{(1)}{\Rightarrow} \text{ExtOwn}(h, Fr_2, \Gamma^r_2(\text{this})) \stackrel{(2)}{\Rightarrow} \\ &\Rightarrow h; FrSt_1 \vdash_{Inv} \ell \Rightarrow h \vdash_{Inv} FrSt_1 \end{aligned}$$

1

Lemma 7.42.

$$\left. \begin{array}{c} h \vdash_{Inv} FrSt \\ \langle t_{new}, h' \rangle = new(h, C, ow) \end{array} \right\} \stackrel{(1)}{(2)} \Rightarrow h \vdash_{Inv} FrSt'$$

Proof.

$$(2) \Rightarrow \forall \iota \in \text{dom}(\mathbf{h}') \setminus \{\iota_{new}\} : \text{owner}(\mathbf{h}, \iota) = \text{owner}(\mathbf{h}', \iota) \stackrel{(1)}{\Rightarrow} \mathbf{h}' \vdash_{Inv} \text{FrSt}$$

□

Lemma 7.43.

$$\left. \begin{array}{l} \text{Fr} = \langle \Gamma^r, \Gamma, U \rangle \\ \text{FrSt} = \text{Fr} \circ \text{FrSt}_0 \\ \Gamma(y) = \text{rep} \langle Cl \rangle \\ \mathbf{h} \vdash_{wf} \text{FrSt} \\ \text{FrSt} \vdash_{Inv} \mathbf{h} \\ \text{ow} = \langle \text{null}^r, \text{newClt}(\mathbf{h}) \rangle \\ \mathbf{h}' = \text{transferCl}(\mathbf{h}, \Gamma^r(y), \text{ow}) \end{array} \right\} \begin{array}{l} (1) \\ (2) \\ (3) \\ (4) \\ (5) \end{array} \Rightarrow \mathbf{h}' \vdash_{Inv} \text{FrSt}$$

Proof.

$$\begin{aligned} \forall Cl \in \text{CltRep} \iota \in \text{dom}(\mathbf{h})(4) \wedge (5) \Rightarrow \text{RefNum}(\mathbf{h}', \text{FrSt}, \langle Cl, \iota \rangle, \text{rep} \langle Cl \rangle) &\stackrel{(1) \wedge (2) \wedge (3)}{\leq} \\ &\leq \text{RefNum}(\mathbf{h}, \text{FrSt}, \langle Cl, \text{null}^r \rangle, \text{uniq}) \leq 1 \Rightarrow \mathbf{h}' \vdash_{Inv} \text{FrSt} \end{aligned}$$

□

Lemma 7.44.

$$\left. \begin{array}{l} \text{Fr} = \langle \Gamma^r, \Gamma, U \rangle \\ \text{FrSt} = \text{Fr} \circ \text{FrSt}_0 \\ \Gamma(y) = \text{uniq} \\ \mathbf{h} \vdash_{wf} \text{FrSt} \\ \text{FrSt} \vdash_{Inv} \mathbf{h} \\ \text{ThisOw}(\mathbf{h}, \text{FrSt}) \\ \mathbf{h}' = \text{transferCl}(\mathbf{h}, \Gamma^r(y), \text{ow}) \end{array} \right\} \begin{array}{l} (1) \\ (2) \\ (3) \\ (4) \\ (5) \end{array} \Rightarrow \mathbf{h}' \vdash_{Inv} \text{FrSt}$$

Proof.

$$\begin{aligned} \forall Cl \in \text{CltRep} \iota \in \text{dom}(\mathbf{h})(4) \wedge (5) \Rightarrow \text{RefNum}(\mathbf{h}, \text{FrSt}, \langle Cl, \iota \rangle, \text{rep} \langle Cl \rangle) = 0 &\stackrel{(1) \wedge (2) \wedge (3)}{\Rightarrow} \\ &\leq \text{RefNum}(\mathbf{h}'L, \text{FrSt}, \langle Cl, \text{null}^r \rangle, \text{uniq}) = 0 \Rightarrow \mathbf{h}' \vdash_{Inv} \text{FrSt} \end{aligned}$$

□

7.8 Ownership tree's lemmas

Lemma 7.45.

$$\left. \begin{array}{l} \text{WfTree}(h) \\ \iota_0 \in \text{dom}(h) \\ h' = h[\iota_0 \cdot f := \iota] \end{array} \right\} \stackrel{(1)}{\quad} \stackrel{(2)}{\quad} \stackrel{(3)}{\quad} \Rightarrow \text{WfTree}(h')$$

Proof.

$$\begin{aligned} \forall \iota \in \text{dom}(h') & \stackrel{(1)}{\Rightarrow} \iota \notin \text{owners}(h, \iota) \downarrow_2 \stackrel{(2) \wedge (3)}{\Rightarrow} \\ & \Rightarrow \iota \notin \text{owners}(h', \iota) \downarrow_2 \Rightarrow \text{WfTree}(h') \end{aligned}$$

□

Lemma 7.46.

$$\left. \begin{array}{l} \text{WfTree}(h) \\ \langle \iota_{new}, h' \rangle = \text{new}(h, C, ow) \end{array} \right\} \stackrel{(1)}{\quad} \stackrel{(2)}{\quad} \Rightarrow \text{WfTree}(h')$$

Proof.

$$\begin{aligned} \forall \iota \in \text{dom}(h') \setminus \{\iota_{new}\} & \stackrel{(1)}{\Rightarrow} \iota \notin \text{owners}(h, \iota) \downarrow_2 \stackrel{(2)}{\Rightarrow} \\ & \Rightarrow \iota \notin \text{owners}(h', \iota) \downarrow_2 \Rightarrow \text{WfTree}(h') \end{aligned}$$

□

Lemma 7.47.

$$\left. \begin{array}{l} \text{WfTree}(h) \\ h' = \text{transferCl}(h, \iota_{tr}, ow) \\ ow = \langle \text{newClt}(h), \text{null}^r \rangle \end{array} \right\} \stackrel{(1)}{\quad} \stackrel{(2)}{\quad} \stackrel{(3)}{\quad} \Rightarrow \text{WfTree}(h')$$

Proof.

$$\begin{aligned} \forall \iota \in \text{dom}(h') \setminus \{\iota_{new}\} & \stackrel{(1)}{\Rightarrow} \iota \notin \text{owners}(h, \iota) \downarrow_2 \stackrel{(2) \wedge (3)}{\Rightarrow} \\ & \Rightarrow \iota \notin \text{owners}(h', \iota) \downarrow_2 \Rightarrow \text{WfTree}(h') \end{aligned}$$

□

Lemma 7.48.

$$\left. \begin{array}{l} \text{WfTree}(h) \\ h' = \text{transferCl}(h, \iota_{tr}, ow) \\ ow = \text{mod20w}(h, m, \Gamma^r(\text{this})) \\ m \in \{\text{rep } \langle Cl \rangle, \text{ peer}\} \\ \text{ThisOw}(h, FrSt) \\ h; _ \vdash_{wf} \iota_{tr} : \text{uniq} \end{array} \right\} \stackrel{(1)}{\quad} \stackrel{(2)}{\quad} \stackrel{(3)}{\quad} \stackrel{(4)}{\quad} \stackrel{(5)}{\quad} \stackrel{(6)}{\quad} \Rightarrow \text{WfTree}(h')$$

Proof.

$$\begin{aligned} \forall \iota \in \text{dom}(h') \setminus \{\iota_{\text{new}}\} & \stackrel{(1)}{\Rightarrow} \iota \notin \text{owners}(h, \iota) \downarrow_2 \stackrel{(2) \wedge (5) \wedge (6)}{\Rightarrow} \\ & \Rightarrow \iota \notin \text{owners}(h', \iota) \downarrow_2 \Rightarrow \text{WfTree}(h') \end{aligned}$$

□

7.9 This owners's lemmas

Lemma 7.49.

$$\left. \begin{array}{l} \text{FrSt} = \langle \Gamma^r_0, \Gamma_0, U_0 \rangle \circ \text{FrSt}_0 \\ \text{FrSt}' = \langle \Gamma'^r_0, \Gamma_0, U'_0 \rangle \circ \text{FrSt}_0 \\ \text{ThisOw}(h, \text{FrSt}) \\ \Gamma^r(\text{this}) = \Gamma'^r(\text{this}) \end{array} \right\} \stackrel{(1)}{\Rightarrow} \text{ThisOw}(h, \text{FrSt}')$$

Proof.

$$\begin{aligned} \text{FrSt}_0 &= \overline{\langle \Gamma^r, \Gamma, U \rangle} \\ (1) \Rightarrow & \left\{ \begin{array}{l} \forall Cl \in \text{CltFree} : \langle Cl, \text{null}^r \rangle \notin \text{owners}(h, \Gamma^r_0(\text{this})) \quad (3) \\ \text{owner}(h, \Gamma^r(\text{this})) \in \text{owners}(h, \Gamma^r_0(\text{this})) \quad (4) \end{array} \right. \\ (3) \wedge (1) \wedge (2) &\Rightarrow \forall Cl \in \text{CltFree} : \langle Cl, \text{null}^r \rangle \notin \text{owners}(h, \Gamma'^r_0(\text{this})) \quad (5) \\ (4) \wedge (1) \wedge (2) &\Rightarrow \overline{\text{owner}(h, \Gamma^r(\text{this})) \in \text{owners}(h, \Gamma'^r_0(\text{this}))} \quad (6) \\ (5) \wedge (6) &\Rightarrow \text{ThisOw}(h, \text{FrSt}') \end{aligned}$$

□

Lemma 7.50.

$$\left. \begin{array}{l} \text{FrSt} = \langle \Gamma^r_0, \Gamma_0, U_0 \rangle \circ \text{FrSt}_0 \\ \text{FrSt}' = \langle \Gamma'^r_0, \Gamma_0, U'_0 \rangle \circ \text{FrSt}_0 \\ \text{ThisOw}(h, \text{FrSt}) \\ \iota_0 \in \text{dom}(h) \\ h' = h[\iota_0 . f := \iota] \end{array} \right\} \stackrel{(1)}{\Rightarrow} \text{ThisOw}(h, \text{FrSt}')$$

Proof.

$$\begin{aligned} \text{FrSt}_0 &= \overline{\langle \Gamma^r, \Gamma, U \rangle} \\ (1) \Rightarrow & \left\{ \begin{array}{l} \forall Cl \in \text{CltFree} : \langle Cl, \text{null}^r \rangle \notin \text{owners}(h, \Gamma^r_0(\text{this})) \quad (4) \\ \text{owner}(h, \Gamma^r(\text{this})) \in \text{owners}(h, \Gamma^r_0(\text{this})) \quad (5) \end{array} \right. \\ (4) \wedge (1) \wedge (3) &\Rightarrow \forall Cl \in \text{CltFree} : \langle Cl, \text{null}^r \rangle \notin \text{owners}(h, \Gamma'^r_0(\text{this})) \quad (6) \\ (5) \wedge (1) \wedge (3) &\Rightarrow \overline{\text{owner}(h, \Gamma^r(\text{this})) \in \text{owners}(h, \Gamma'^r_0(\text{this}))} \quad (7) \\ (6) \wedge (7) &\Rightarrow \text{ThisOw}(h, \text{FrSt}') \end{aligned}$$

□

Lemma 7.51.

$$\left. \begin{array}{l} Fr = \langle \Gamma^r_0, \Gamma_0, U_0 \rangle \\ Fr' = \langle \Gamma'^r_0, \Gamma'_0, U'_0 \rangle \\ FrSt = Fr \circ FrSt_0 \\ FrSt' = Fr' \circ FrSt \\ ThisOw(h, FrSt) \\ h; \Gamma^r(this) \vdash_{wf} \Gamma'^r(this) : m \\ IsWritable(m) \end{array} \right\} \stackrel{(1)}{\Rightarrow} ThisOw(h, FrSt')$$

Proof.

$$\begin{aligned} FrSt_0 &= \overline{\langle \Gamma^r, \Gamma, U \rangle} \\ (1) &\Rightarrow \left\{ \begin{array}{l} \forall Cl \in \text{CltFree} : \langle Cl, \text{null}^r \rangle \notin \text{owners}(h, \Gamma^r_0(this)) \quad (4) \\ \text{owner}(h, \Gamma^r(this)) \in \text{owners}(h, \Gamma^r_0(this)) \end{array} \right. \quad (5) \\ (4) \wedge (1) \wedge (3) &\Rightarrow \forall Cl \in \text{CltFree} : \langle Cl, \text{null}^r \rangle \notin \text{owners}(h, \Gamma'^r_0(this)) \quad (6) \\ (5) \wedge (1) \wedge (3) &\Rightarrow \text{owner}(h, \Gamma^r(this)) \in \text{owners}(h, \Gamma'^r_0(this)) \quad (7) \\ (1) \wedge (3) &\Rightarrow \text{owner}(h, \Gamma^r_0(this)) \in \text{owners}(h, \Gamma'^r_0(this)) \quad (8) \\ (6) \wedge (7) \wedge (8) &\Rightarrow ThisOw(h, FrSt') \end{aligned}$$

□

Lemma 7.52.

$$\left. \begin{array}{l} ThisOw(h, FrSt) \\ \langle \iota, h' \rangle = new(h, C, ow) \end{array} \right\} \stackrel{(1)}{\Rightarrow} \stackrel{(2)}{\Rightarrow} ThisOw(h', FrSt)$$

Proof.

$$\begin{aligned} FrSt_0 &= \overline{\langle \Gamma^r, \Gamma, U \rangle} \\ (1) &\Rightarrow \left\{ \begin{array}{l} \forall Cl \in \text{CltFree} : \langle Cl, \text{null}^r \rangle \notin \text{owners}(h, \Gamma^r_0(this)) \quad (3) \\ \text{owner}(h, \Gamma^r(this)) \in \text{owners}(h, \Gamma^r_0(this)) \end{array} \right. \quad (4) \\ (3) \wedge (1) \wedge (2) &\Rightarrow \forall Cl \in \text{CltFree} : \langle Cl, \text{null}^r \rangle \notin \text{owners}(h', \Gamma^r_0(this)) \quad (5) \\ (4) \wedge (1) \wedge (2) &\Rightarrow \text{owner}(h', \Gamma^r(this)) \in \text{owners}(h, \Gamma^r_0(this)) \quad (6) \\ (5) \wedge (6) &\Rightarrow ThisOw(h', FrSt) \end{aligned}$$

□

Lemma 7.53.

$$\left. \begin{array}{l} Fr_1 = \langle \Gamma^r_1, \Gamma_1, U_1 \rangle \\ Fr_2 = \langle \Gamma^r_2, \Gamma_2, U_2 \rangle \\ FrSt_1 = Fr_1 \circ FrSt_0 \\ FrSt_2 = Fr_2 \circ FrSt_1 \\ h; \Gamma^r_1(this) \vdash_{wf} \Gamma^r_2(this) : m \quad (1) \\ m \in \{\text{peer}, \text{this}, \text{rep } \langle Cl \rangle\} \quad (2) \\ ThisOw(h, FrSt_2) \quad (3) \end{array} \right\} \Rightarrow ThisOw(h, FrSt_1)$$

Proof.

$$\begin{aligned} FrSt_0 &= \overline{\langle \Gamma^r, \Gamma, U \rangle} \\ (1) \Rightarrow &\left\{ \begin{array}{l} \forall Cl \in \text{CltFree} : \langle Cl, \text{null}^r \rangle \notin \text{owners}(h, \Gamma^r_2(this)) \quad (4) \\ \text{owner}(h, \Gamma^r(this)) \in \text{owners}(h, \Gamma^r_2(this)) \end{array} \right. \quad (5) \\ (4) \wedge (1) \wedge (2) \Rightarrow &\overline{\forall Cl \in \text{CltFree} : \langle Cl, \text{null}^r \rangle \notin \text{owners}(h, \Gamma^r_1(this))} \quad (6) \\ (5) \wedge (1) \wedge (2) \Rightarrow &\overline{\text{owner}(h, \Gamma^r(this)) \in \text{owners}(h, \Gamma^r_1(this))} \quad (7) \\ (6) \wedge (7) \Rightarrow &ThisOw(h, FrSt_1) \end{aligned}$$

□

Lemma 7.54.

$$\left. \begin{array}{l} Fr = \langle \Gamma^r_0, \Gamma_0, U_0 \rangle \\ ThisOw(h, FrSt) \quad (1) \\ h' = transferCl(h, \iota_{tr}, ow) \quad (2) \\ h; \Gamma^r(this) \vdash_{wf} \iota_{tr} : m \quad (3) \\ m \in \{\text{rep } \langle Cl \rangle, \text{uniqu}\} \quad (4) \end{array} \right\} \Rightarrow ThisOw(h', FrSt)$$

Proof.

$$\begin{aligned} FrSt_0 &= \overline{\langle \Gamma^r, \Gamma, U \rangle} \\ (1) \Rightarrow &\left\{ \begin{array}{l} \forall Cl \in \text{CltFree} : \langle Cl, \text{null}^r \rangle \notin \text{owners}(h, \Gamma^r_0(this)) \quad (5) \\ \text{owner}(h, \Gamma^r(this)) \in \text{owners}(h, \Gamma^r_0(this)) \end{array} \right. \quad (6) \\ (5) \wedge (2) \wedge (4) \Rightarrow &\overline{\forall Cl \in \text{CltFree} : \langle Cl, \text{null}^r \rangle \notin \text{owners}(h', \Gamma^r_0(this))} \quad (7) \\ (6) \wedge (2) \wedge (4) \Rightarrow &\overline{\text{owner}(h', \Gamma^r(this)) \in \text{owners}(h', \Gamma^r_0(this))} \quad (8) \\ (7) \wedge (8) \Rightarrow &ThisOw(h', FrSt) \end{aligned}$$

□

7.10 Owner as modifier's lemmas

Lemma 7.55.

$$\left. \begin{array}{l} Fr = \langle \Gamma^r, \Gamma, U \rangle \\ Fr' = \langle \Gamma^r, \Gamma', U' \rangle \\ \{v \in U' | \Gamma'(v) = \text{uniq}\} \subseteq \{v \in U | \Gamma(v) = \text{uniq}\} \quad (1) \\ \text{owner}(h, \Gamma^r(this)) \in \text{owners}(h, \Gamma'^r(this)) \quad (2) \end{array} \right\} \Rightarrow \text{OwAsMod}(h, h, Fr, Fr')$$

Proof.

$$\begin{aligned} \forall \iota \in \text{dom}(h) : \neg \text{ExtOwn}(h, Fr, \iota) &\stackrel{(1)\wedge(2)}{\Rightarrow} \\ \Rightarrow h(\iota) = h(\iota) \wedge \neg \text{ExtOwn}(h, Fr', \iota) &\Rightarrow \text{OwAsMod}(h, h, Fr, Fr') \end{aligned}$$

□

Lemma 7.56.

$$\left. \begin{array}{l} Fr = \langle \Gamma^r, \Gamma, U \rangle \\ \iota_0 \in \text{dom}(h) \quad (1) \\ h, \Gamma^r(this) \vdash_{wf} \iota : m \quad (2) \\ \text{IsWritable}(m) \quad (3) \\ h' = h[\iota_0 . f := \iota] \quad (4) \end{array} \right\} \Rightarrow \text{OwAsMod}(h, h', Fr, Fr)$$

Proof.

$$\begin{aligned} \forall \iota \in \text{dom}(h) : \neg \text{ExtOwn}(h, Fr, \iota) &\stackrel{(3)\wedge(4)}{\Rightarrow} \\ \Rightarrow h(\iota) = h'(\iota) \wedge \neg \text{ExtOwn}(h', Fr, \iota) &\Rightarrow \text{OwAsMod}(h, h', Fr, Fr) \end{aligned}$$

□

Lemma 7.57.

$$\text{OwAsMod}(h, h', Fr, Fr') \quad (1) \Rightarrow \text{OwAsMod}(h, h', Fr, Fr)$$

Proof.

$$\begin{aligned} \forall \iota \in \text{dom}(h) : \neg \text{ExtOwn}(h, Fr, \iota) &\stackrel{(1)}{\Rightarrow} \\ \Rightarrow h(\iota) = h'(\iota) \wedge \neg \text{ExtOwn}(h, Fr', \iota) &\Rightarrow \\ \Rightarrow h(\iota) = h'(\iota) \wedge \neg \text{ExtOwn}(h, Fr, \iota) &\Rightarrow \text{OwAsMod}(h, h', Fr, Fr) \end{aligned}$$

□

Lemma 7.58.

$$\left. \begin{array}{l} Fr = \langle \Gamma^r, \Gamma, U \rangle \\ U' = U \setminus \{x\} \\ \Gamma^{r'} = \Gamma^r[x \mapsto \iota] \\ Fr' = \langle \Gamma^{r'}, \Gamma, U' \rangle \\ \langle \iota, h' \rangle = new(h, C, ow) \end{array} \right\} \Rightarrow OwAsMod(h, h', Fr, Fr')$$

Proof.

$$\begin{aligned} \forall \iota \in \text{dom}(h) : \neg \text{ExtOwn}(h, Fr, \iota) &\stackrel{(1)}{\Rightarrow} \\ \Rightarrow h(\iota) = h'(\iota) \wedge \neg \text{ExtOwn}(h', Fr', \iota) &\Rightarrow OwAsMod(h, h', Fr, Fr') \end{aligned}$$

□

Lemma 7.59.

$$\left. \begin{array}{l} OwAsMod(h, h_1, Fr, Fr_1) \\ OwAsMod(h_1, h_2, Fr_1, Fr_2) \end{array} \right\} \stackrel{(1)}{\Rightarrow} \left. \begin{array}{l} OwAsMod(h, h_2, Fr, Fr_2) \end{array} \right\} \stackrel{(2)}{\Rightarrow} OwAsMod(h, h_2, Fr, Fr_2)$$

Proof.

$$\begin{aligned} \forall \iota \in \text{dom}(h) : \\ \neg \text{ExtOwn}(h, Fr, \iota) &\stackrel{(1)}{\Rightarrow} \left\{ \begin{array}{l} h(\iota) = h_1(\iota) \\ \neg \text{ExtOwn}(h_1, Fr_1, \iota) \end{array} \right. \stackrel{(3)}{\Rightarrow} \\ (4) &\stackrel{(2)}{\Rightarrow} \left\{ \begin{array}{l} h_1(\iota) = h_2(\iota) \\ \neg \text{ExtOwn}(h_2, Fr_2, \iota) \end{array} \right. \stackrel{(5)}{\Rightarrow} \stackrel{(6)}{\Rightarrow} \\ (3) \wedge (5) &\Rightarrow h(\iota) = h_2(\iota) \stackrel{(7)}{\Rightarrow} \\ (6) \wedge (7) &\Rightarrow OwAsMod(h, h_2, Fr, Fr_2) \end{aligned}$$

□

Lemma 7.60.

$$\left. \begin{array}{l} Fr = \langle \Gamma^r, \Gamma, U \rangle \\ U' = U \setminus \{x\} \\ \Gamma^{r'} = \Gamma^r[x \mapsto \Gamma^r(y)] \\ Fr' = \langle \Gamma^{r'}, \Gamma, U' \rangle \\ h; \Gamma^r(this) \vdash_{wf} \Gamma^r(y) : rep \langle Cl \rangle \quad (1) \\ h' = transferCl(h, \Gamma^r(y), ow) \quad (2) \\ \Gamma(x) = uniq \quad (3) \end{array} \right\} \Rightarrow OwAsMod(h, h', Fr, Fr')$$

Proof.

$$\begin{aligned} \forall \iota \in \text{dom}(h) : \neg \text{ExtOwn}(h, Fr, \iota) &\xrightarrow{(1) \wedge (2)} \\ \Rightarrow h(\iota) = h'(\iota) \wedge \neg \text{ExtOwn}(h', Fr', \iota) &\Rightarrow \text{OwAsMod}(h, h', Fr, Fr') \end{aligned}$$

□

Lemma 7.61.

$$\left. \begin{array}{l} Fr = \langle \Gamma^r, \Gamma, U \rangle \\ \Gamma(y) = \text{uniq} \\ h' = \text{transferCl}(h, \Gamma^r(y), \text{ow}) \end{array} \right\} \xrightarrow{(1) \quad (2)} \text{OwAsMod}(h, h', Fr, Fr)$$

Proof.

$$\begin{aligned} \forall \iota \in \text{dom}(h) : \neg \text{ExtOwn}(h, Fr, \iota) &\xrightarrow{(2)} \\ \Rightarrow h(\iota) = h'(\iota) \wedge \neg \text{ExtOwn}(h', Fr', \iota) &\Rightarrow \text{OwAsMod}(h, h', Fr, Fr') \end{aligned}$$

□